

Efficient Formalism- Independent Monitoring of Parametric Properties

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Overview

- Motivation
- Parametric Monitoring
- Enable Sets Based Optimization
- Results
- Conclusion

Monitoring Examples

- Require authentication before allowing access
- Events
 - **authenticate** - when the program authenticates
 - **access** - just before the program accesses the resource
- Property (using past time linear temporal logic)
 - **access** → *eventually in the past authenticate*
- Handler
 - Perform the authenticate

Monitoring Examples

- No write after file close
- Events
 - **open** - the open call for a file
 - **write** - just before a write to a file
 - **close** - the close call for a file
- Property (using a regular expression)
 - **open write*** **close write**
- Handler
 - reopen the file, or disallow the write with a warning

Monitoring Examples

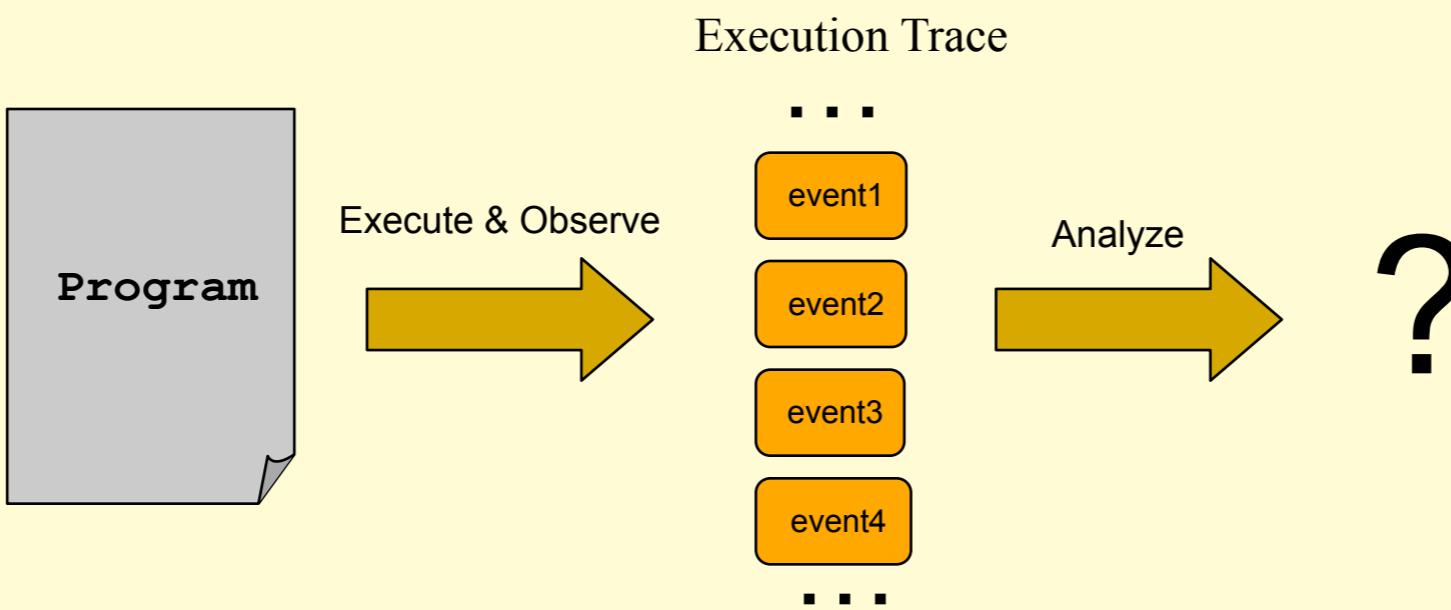
- Releases and acquires of a lock must be matched
- Events
 - **begin** - begin of a method
 - **end** - end of a method
 - **acq** - acquire of a lock
 - **rel** - release of a lock
- Property (using a context free grammar)
 - $S \rightarrow \epsilon \mid S \ S \mid \text{begin } S \ \text{end} \mid \text{acq } S \ \text{rel}$
- Handler
 - Issue an error message

Applications of Monitoring

- Debugging
 - Deployment - development
 - Handler - error messages
- Testing
 - Deployment - development
 - Handler - error messages
- Security/Reliability/Runtime Verification
 - Deployment - production systems
 - Handler - recovery code

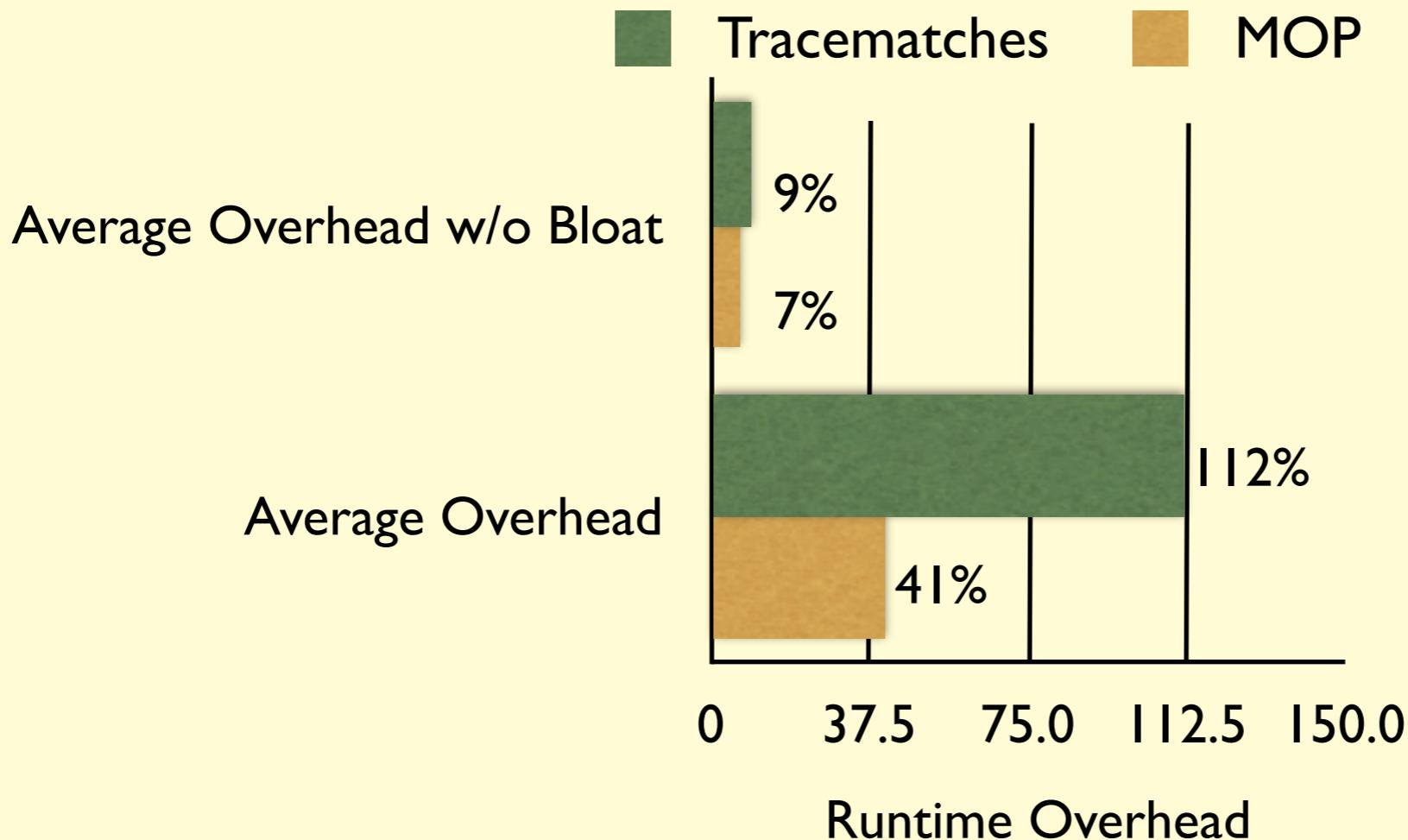
Monitoring in a Nutshell

- Observe run of a system



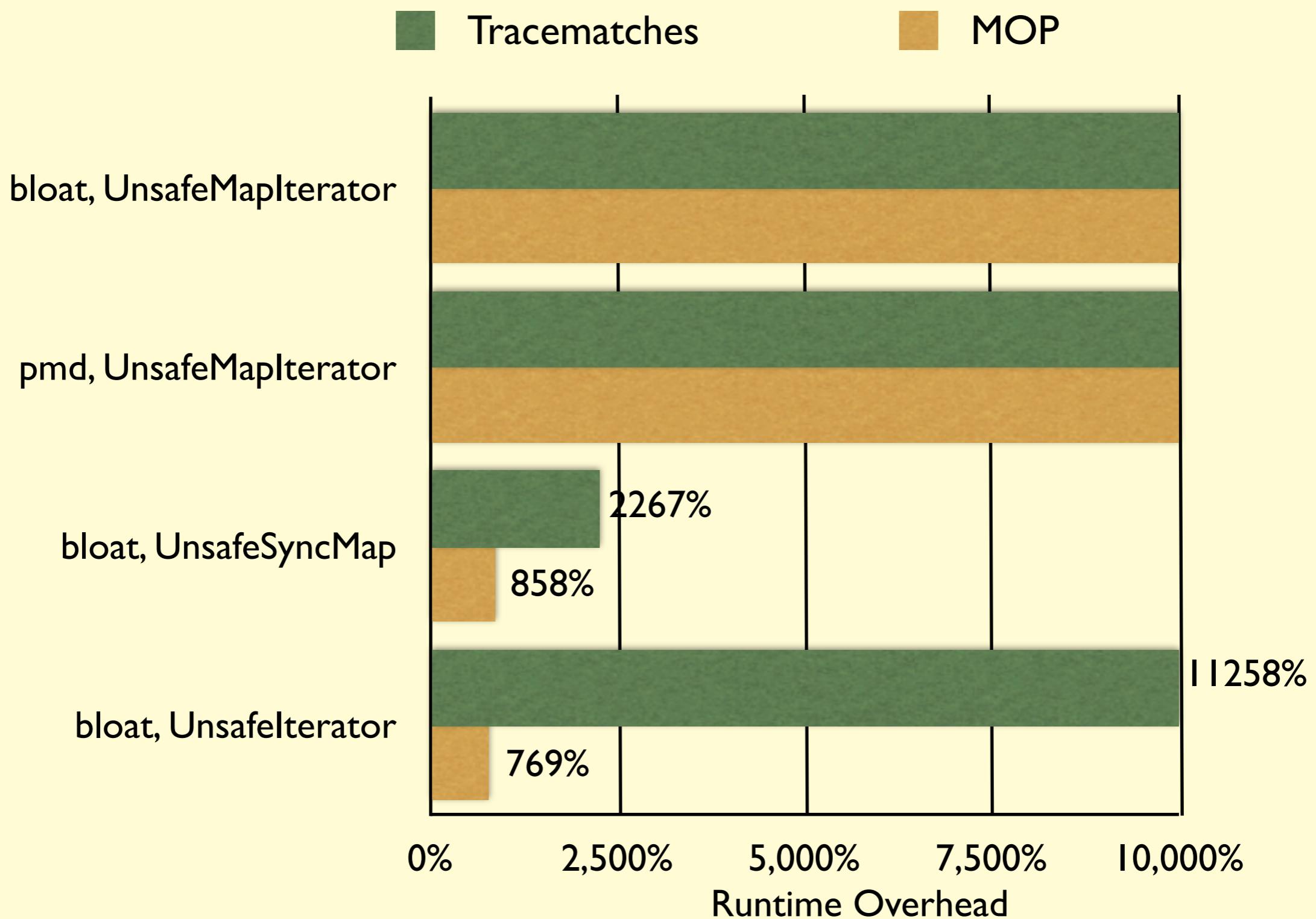
- Analyze it against desired properties
- React/Report using handlers (if needed)
- Scalable!

Expected Overheads

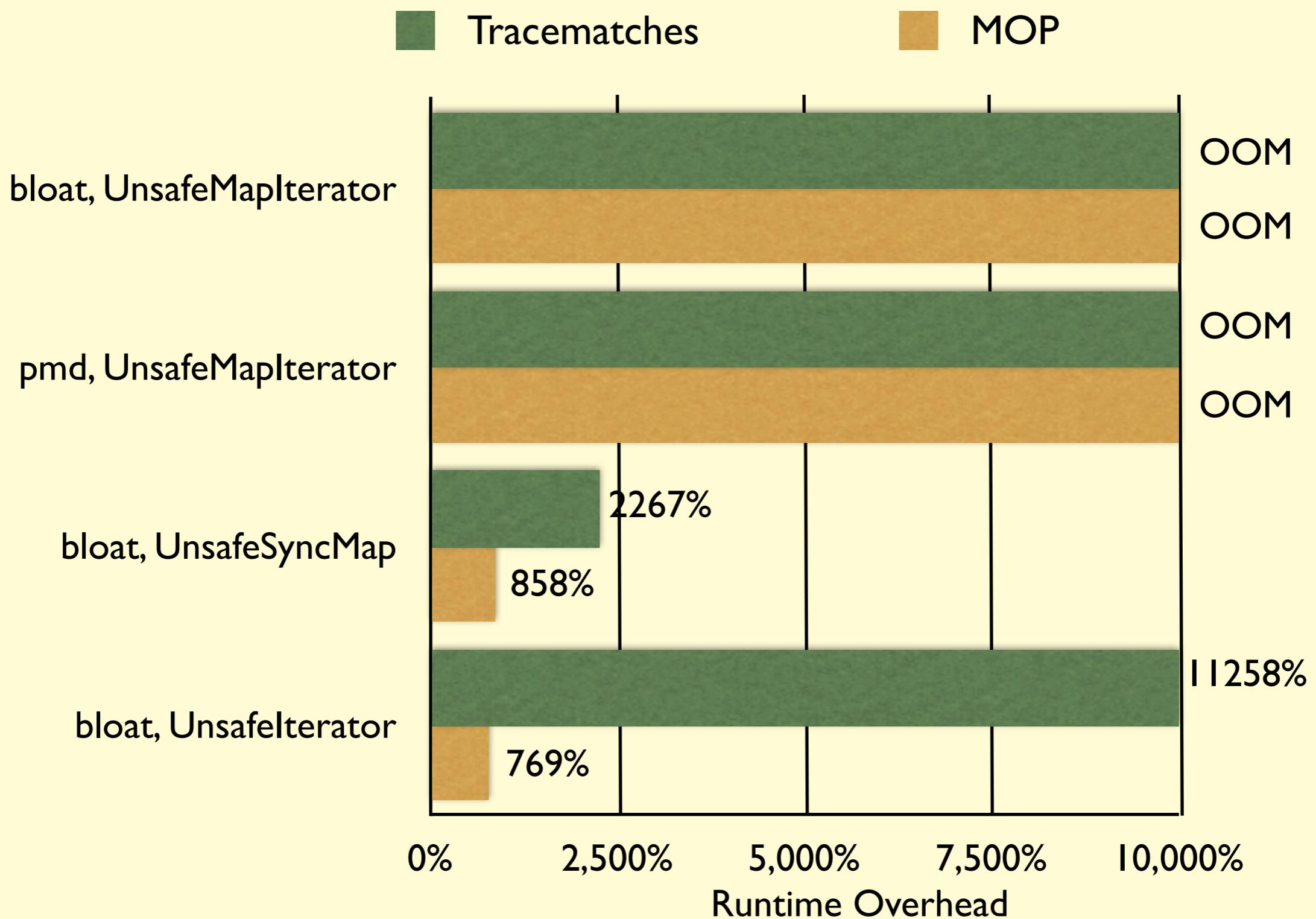


- Tracematches - competing system
- MOP - our system (without optimization)

Bad Overheads



Bad Overheads



Monitoring Systems

- MAC (UPenn)
- PAX (NASA)
- TimeRover (commercial)
- HAWK/Eagle (NASA)
- MOP (UIUC)
- POTA (UTA)
- PQL (Stanford)
- Tracematches (Oxford)
- PTQL (Berkeley/Stanford/Novell)
- Pal (UPenn)

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Parametric Monitoring Example

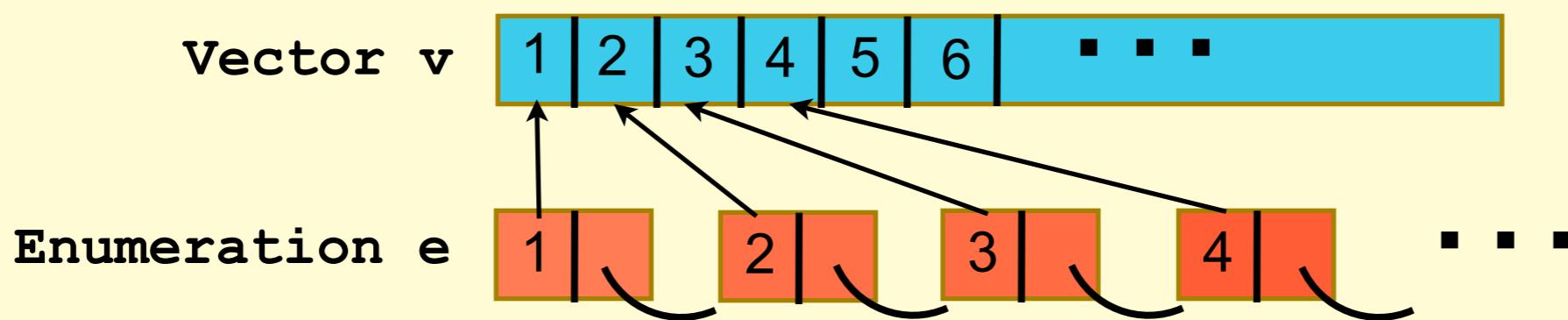
- Actually almost all properties are parametric
 - Events carry parameters
- We wish to apply authenticate before use property to *all* resources in a system
 - **access<r>** → *eventually in the past* **authenticate<r>**
- Allows for multiple resources
 - **authenticate<r>** may require a specific identification per resource

More Parametric Properties

- No write after the close of a specific file
 - `open<f> write<f>* close<f> write<f>`
 - Without parameters one needs to write this property for each file
 - Or close of one file may trigger a handler on write to a *different* file!
- Releases and acquires of a lock must be matched
 - $S \rightarrow \epsilon \mid S \ S \mid \text{begin}<t> \ S \ \text{end}<t> \mid \text{acq}<\text{l}> \ S \ \text{rel}<\text{l}>$
 - We only wish to match acquires and release to the *same* lock
 - The t (thread) parameter can further ensure that the locks are matched within a given thread

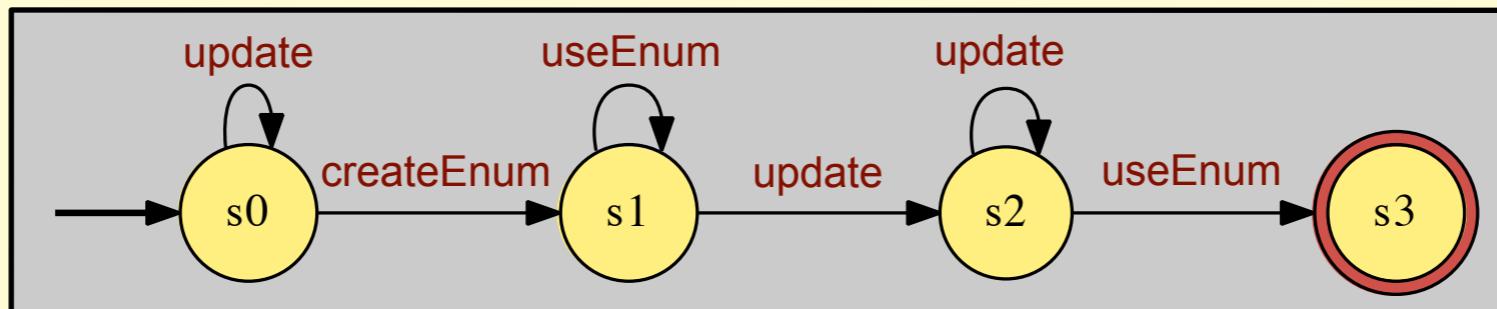
Running Example: Unsafe Enumeration

- One should not change a vector while being accessed via one or more enumeration objects



Unsafe Enumeration as a Parametric Property

- Violation pattern of three events
 - `update<v>` : change in vector `v`
 - `createEnum<v,e>` : create enumeration `e` from vector `v`
 - `useEnum<e>` : use enumeration `e`
- `update* createEnum useEnum* update+ useEnum`



Generic Parametric Monitoring Idea

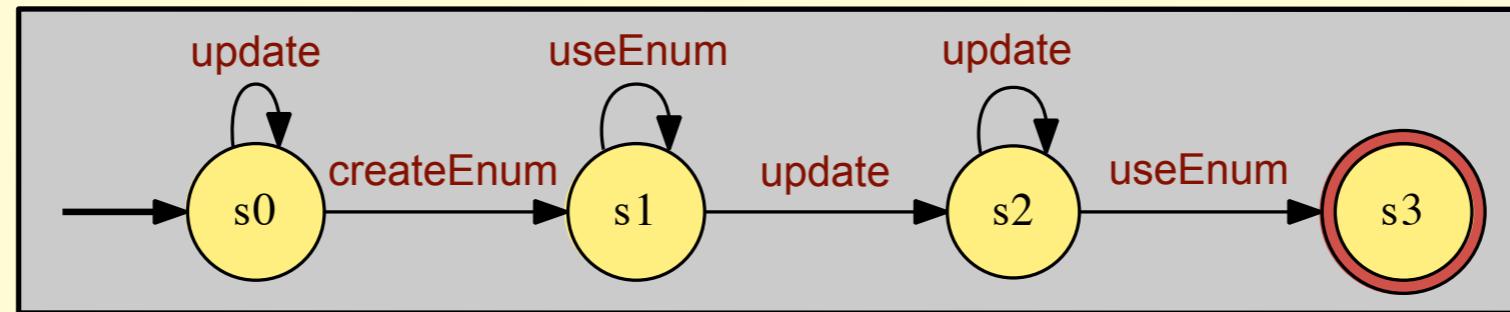
[TACAS 2009]

- Keep one monitor for each *parameter instance*
 - A parameter instance is the bindings of specific objects to its parameter
 - E.g., $\langle v \rightarrow v_2, e \rightarrow e_3 \rangle$
- Each monitor knows nothing of parameters, operates on one trace slice

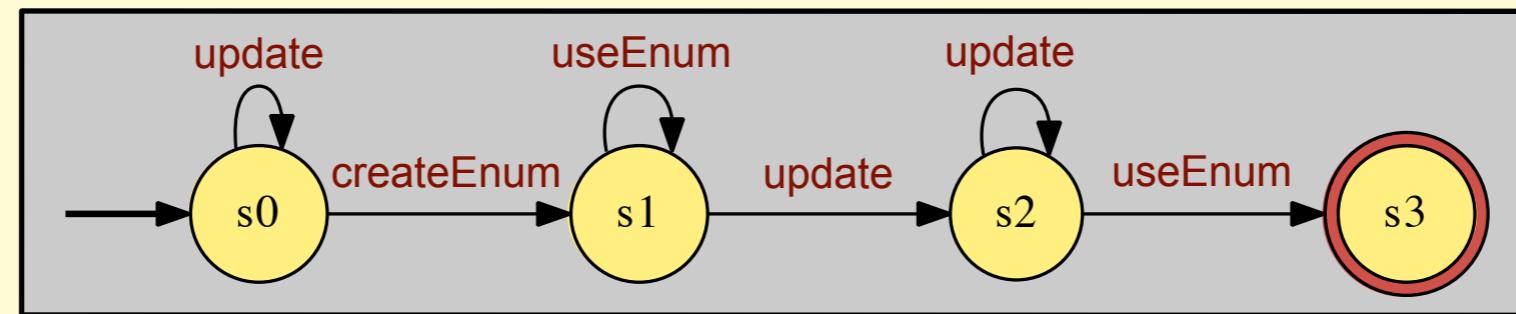
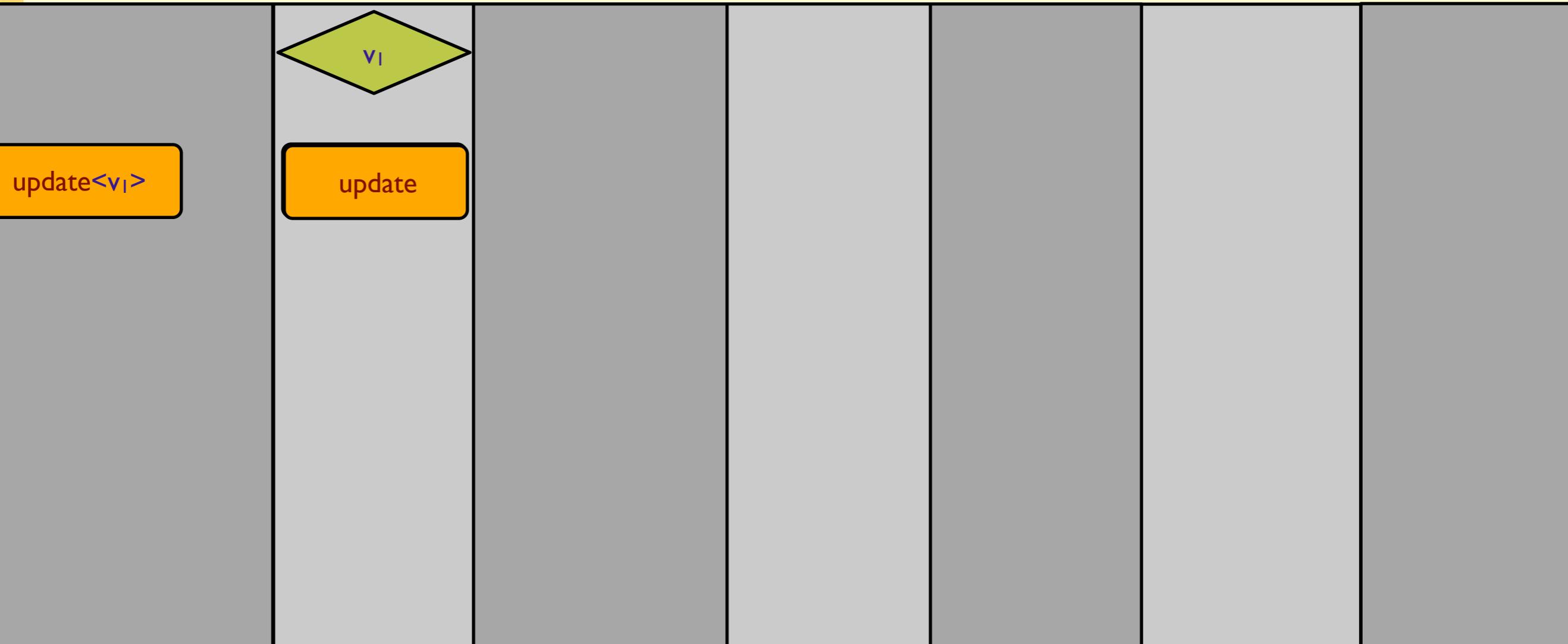
Possible trace - $\text{createEnum} \langle v_1, e_1 \rangle \text{ createEnum} \langle v_1, e_2 \rangle$
 $\text{useEnum} \langle e_1 \rangle \text{ update} \langle v_1 \rangle \text{ useEnum} \langle e_2 \rangle$

$\langle v_1, e_1 \rangle$ slice - createEnum useEnum update
 $\langle v_1, e_2 \rangle$ slice - createEnum update useEnum

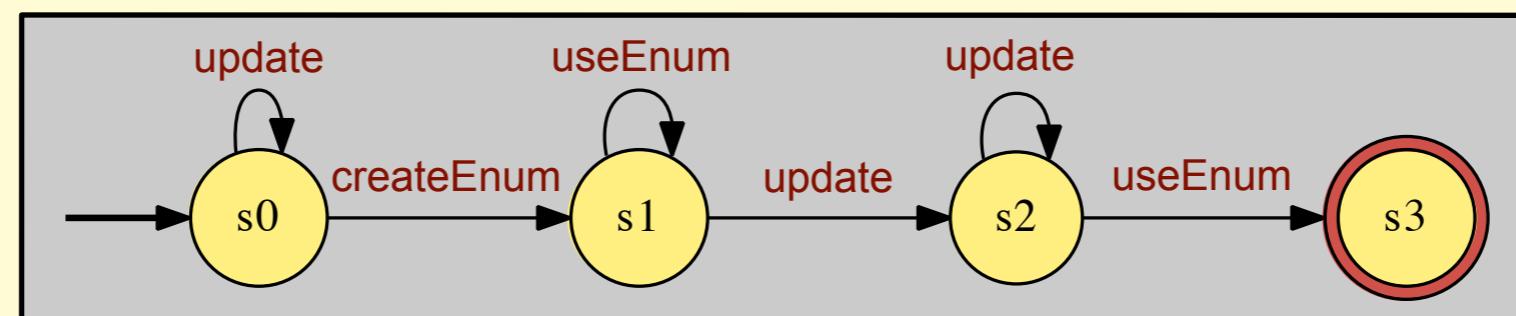
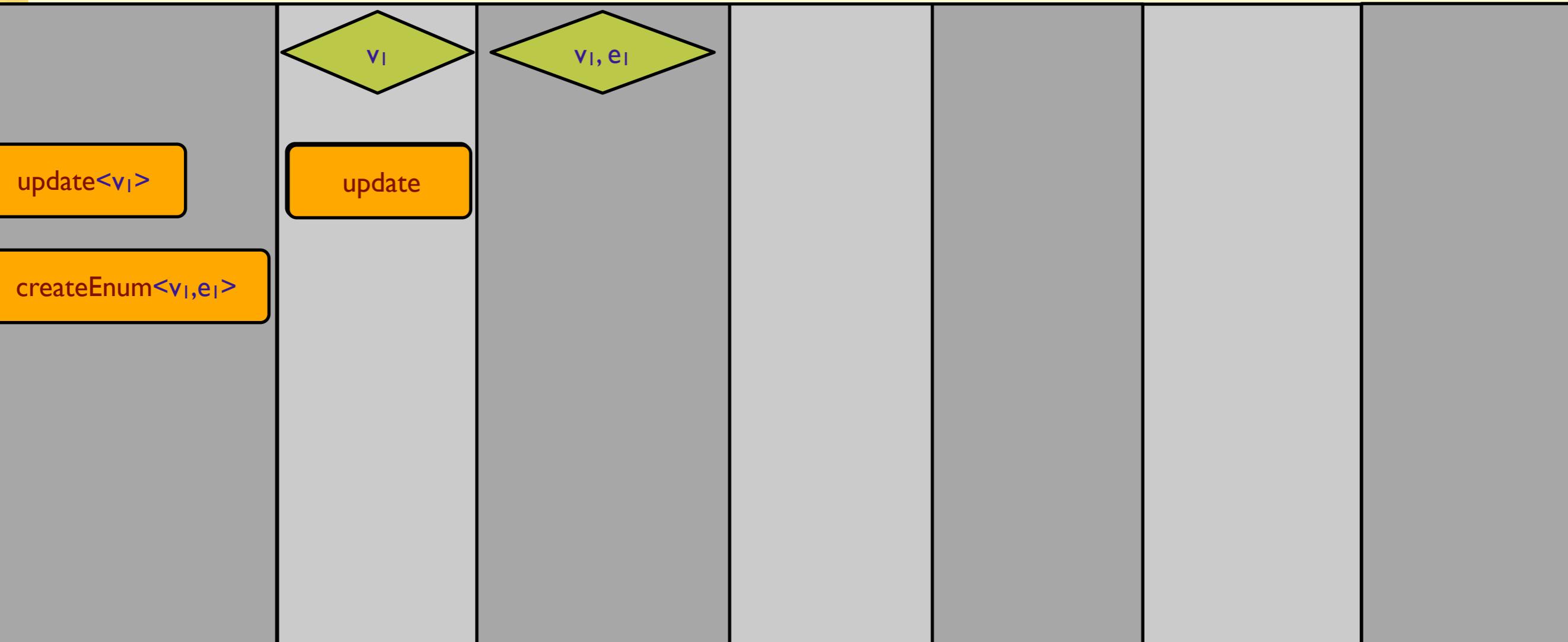
Example Run



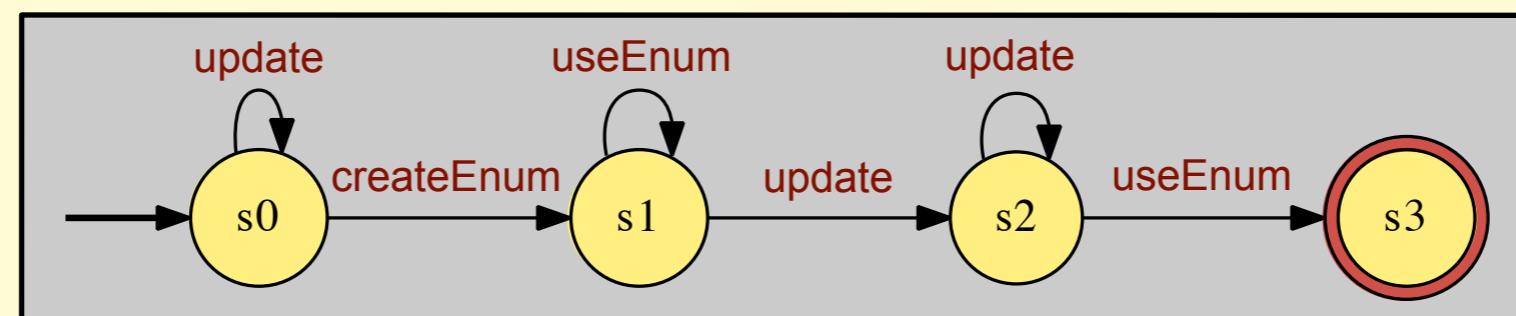
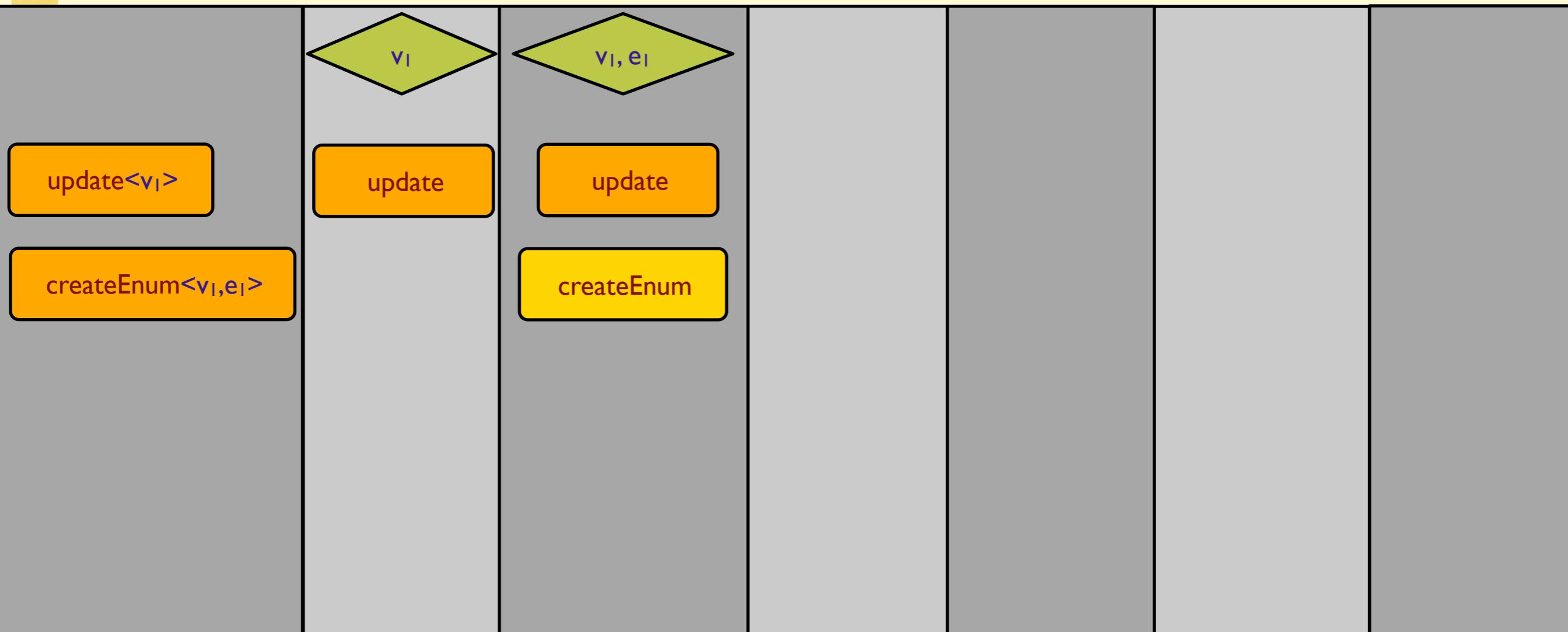
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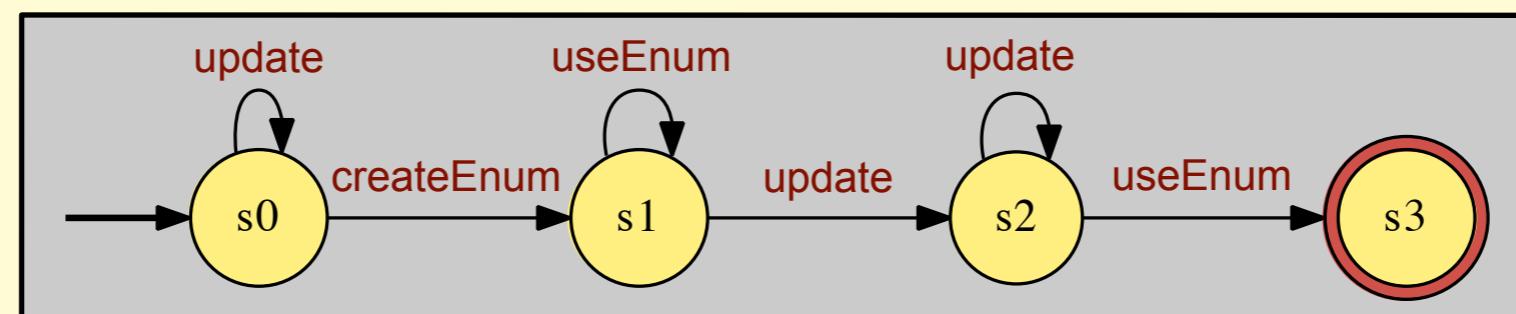
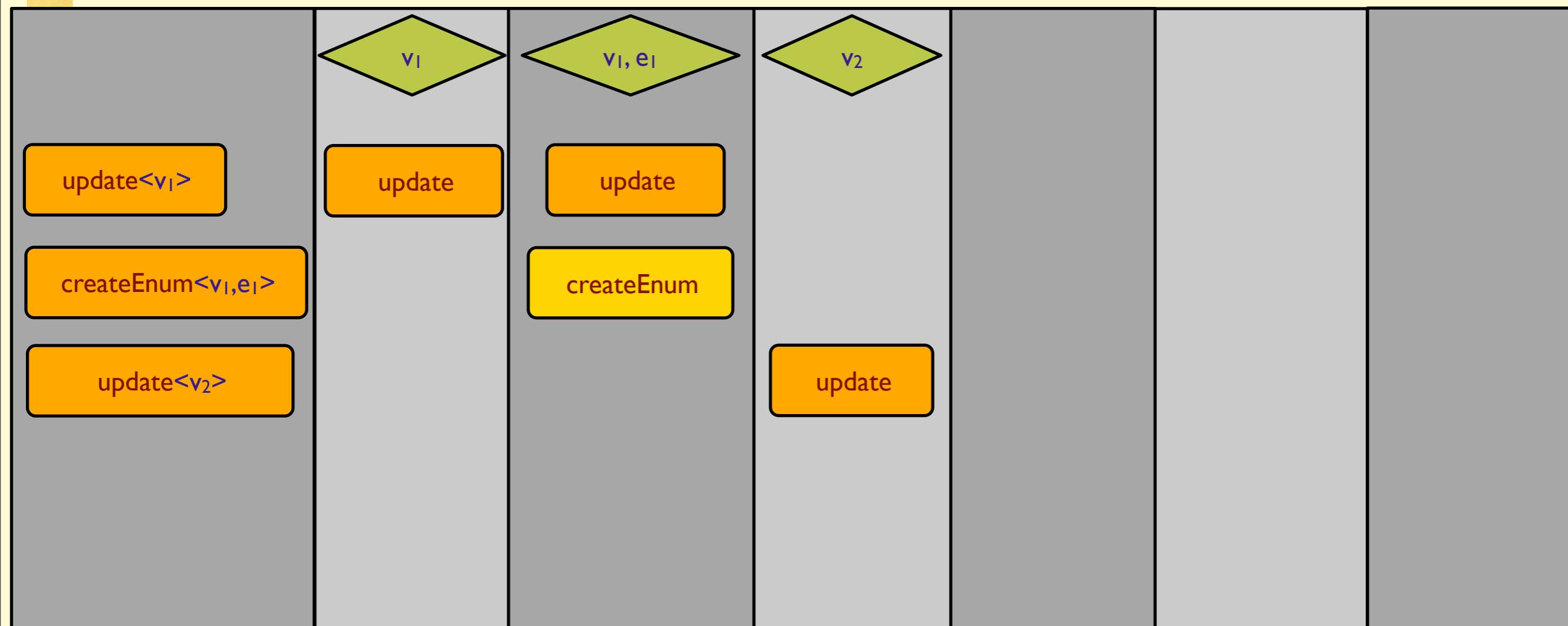
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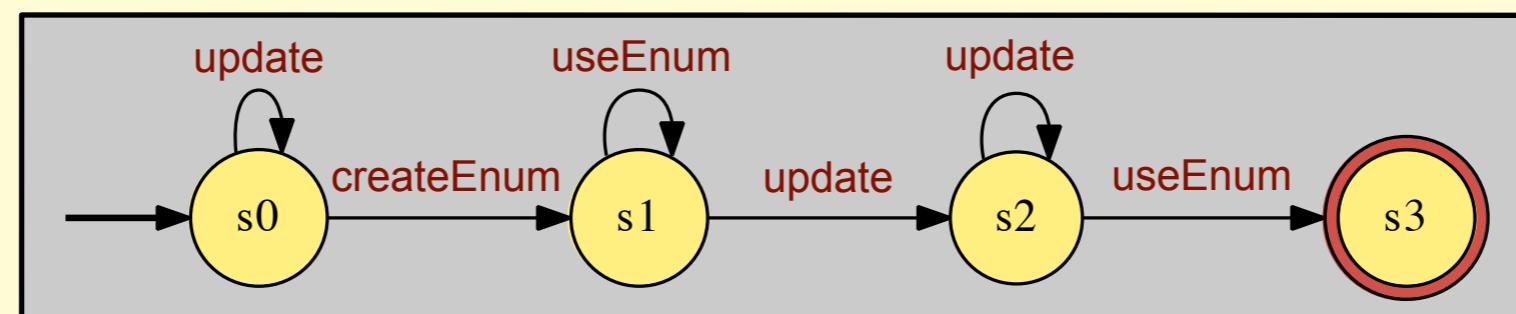
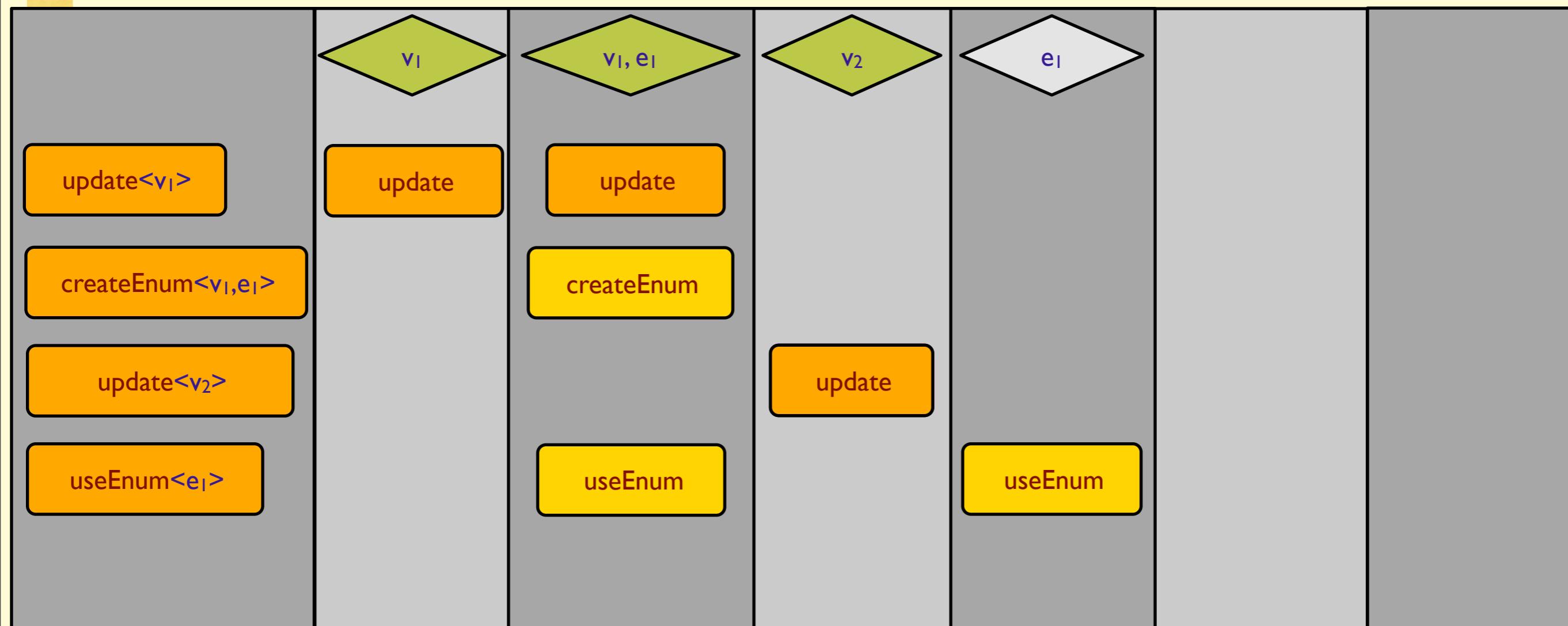
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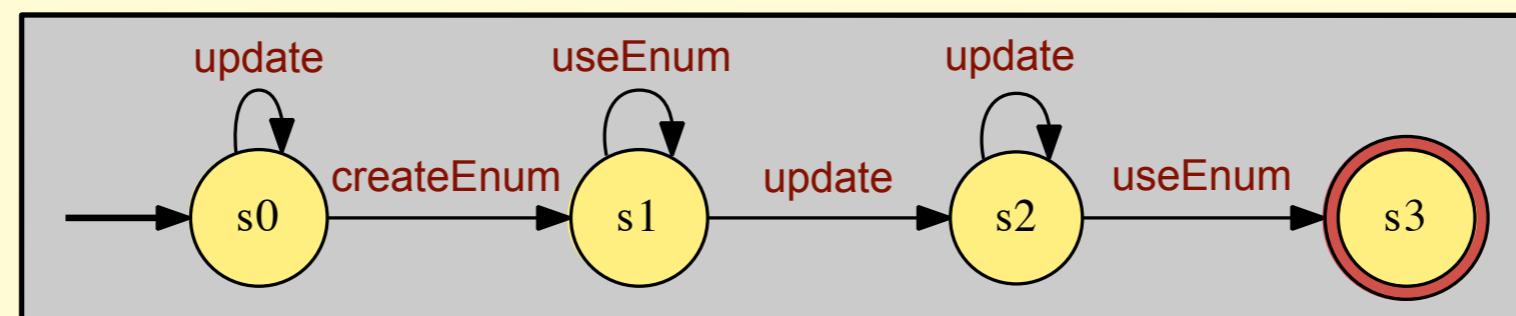
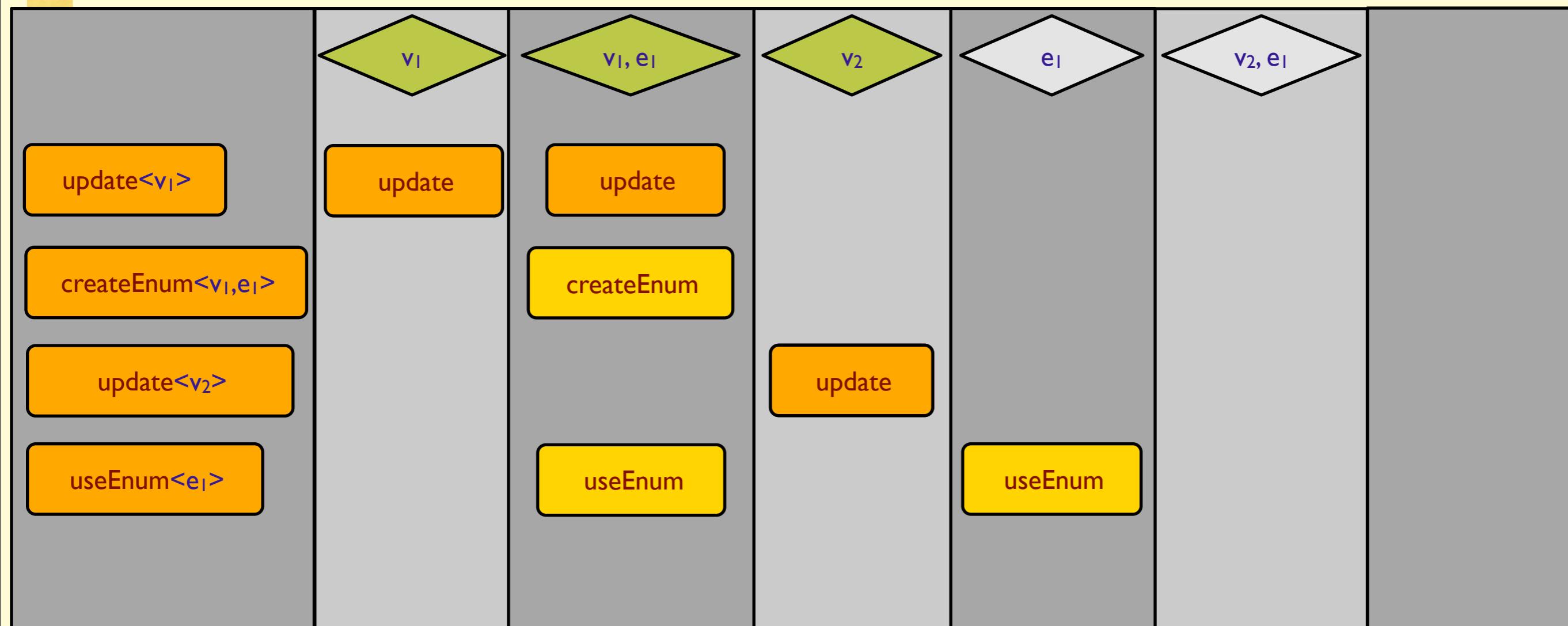
Example Run



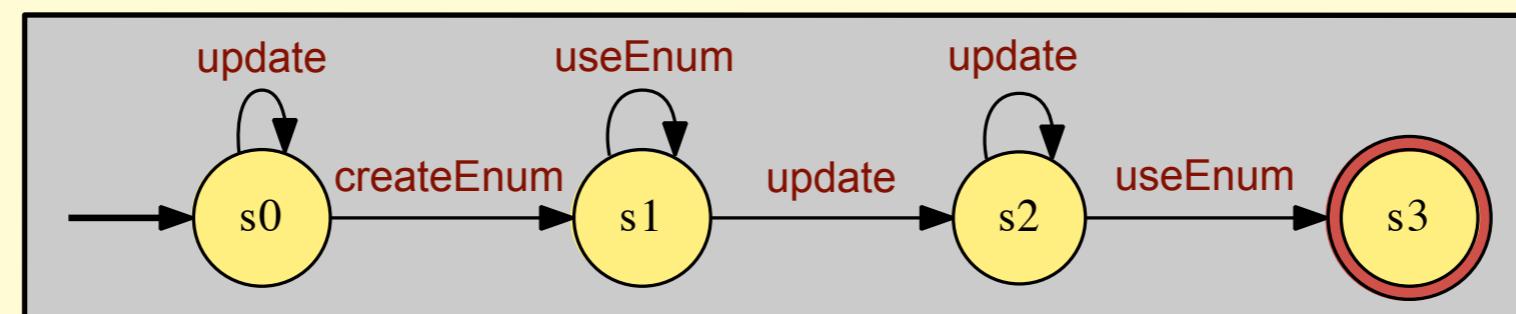
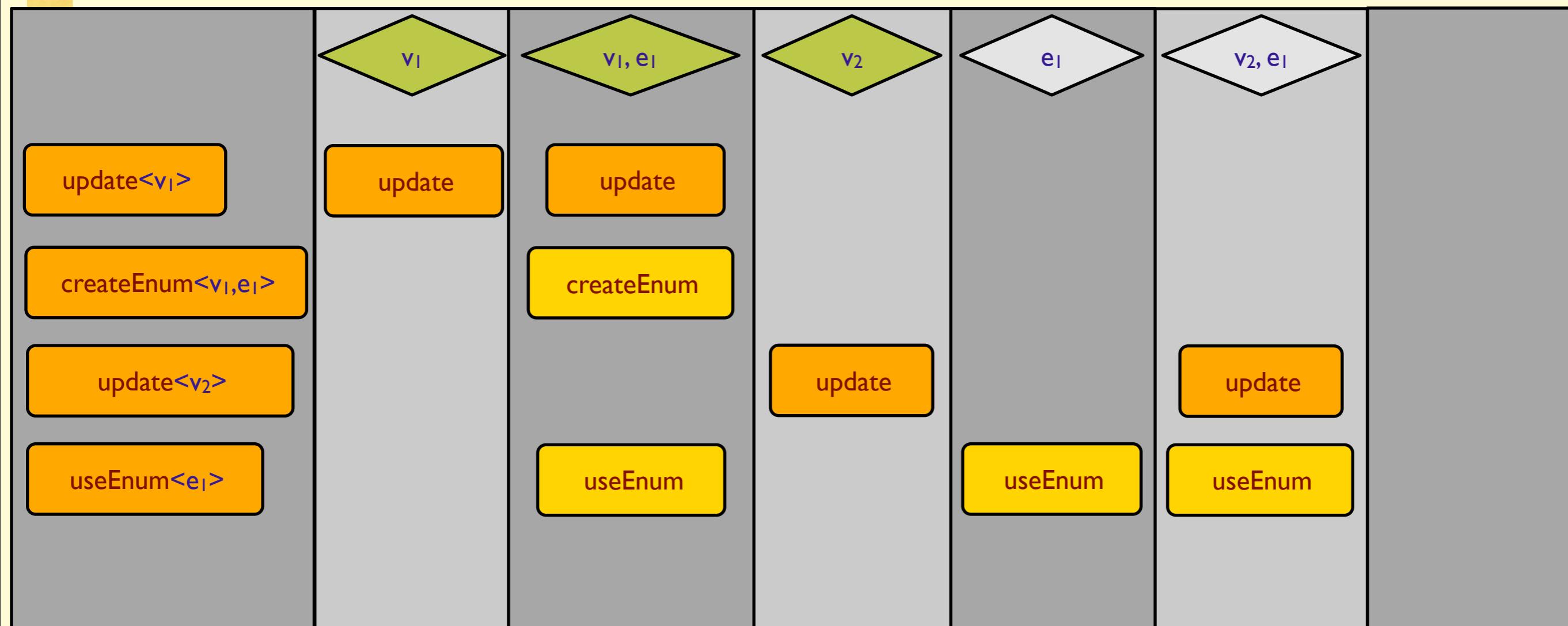
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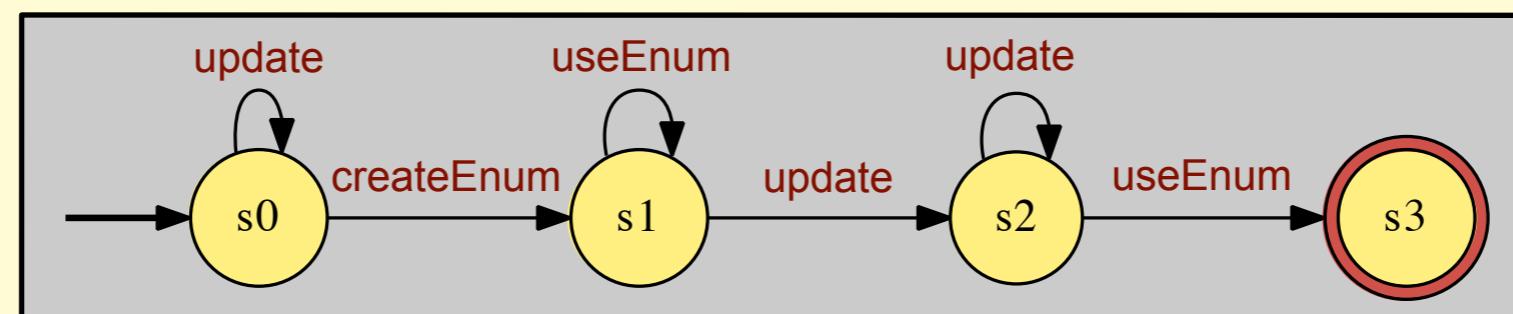
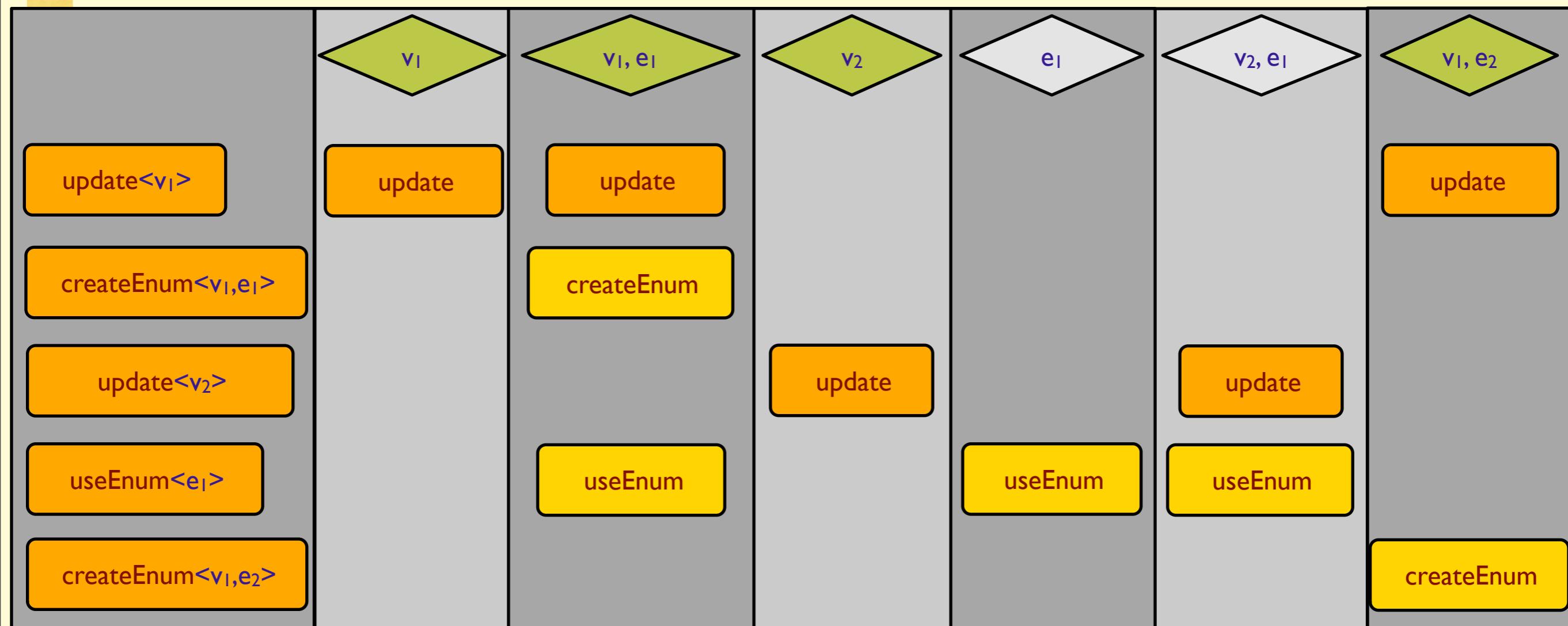
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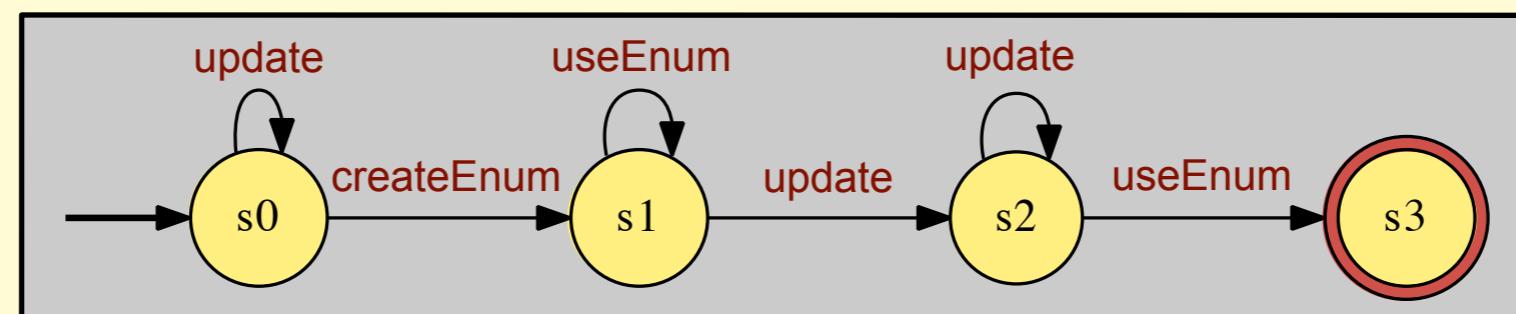
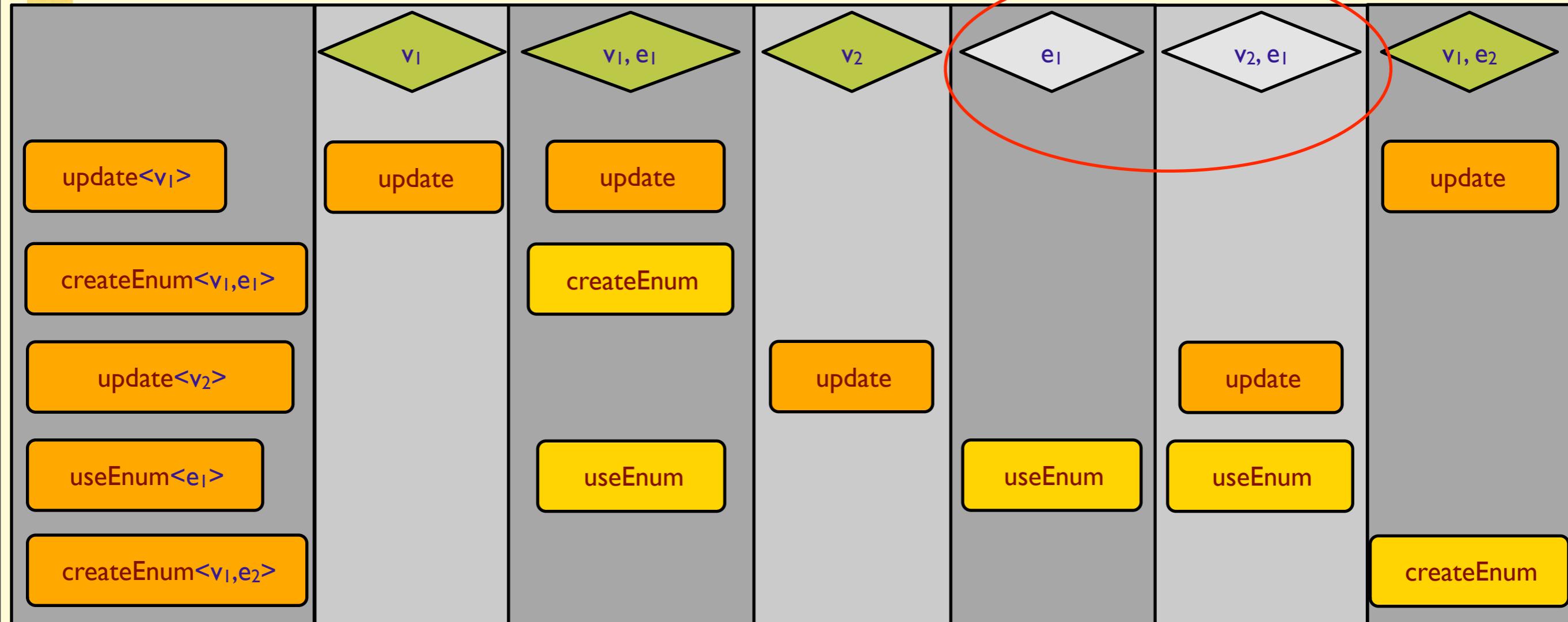
Example Run



Example Run

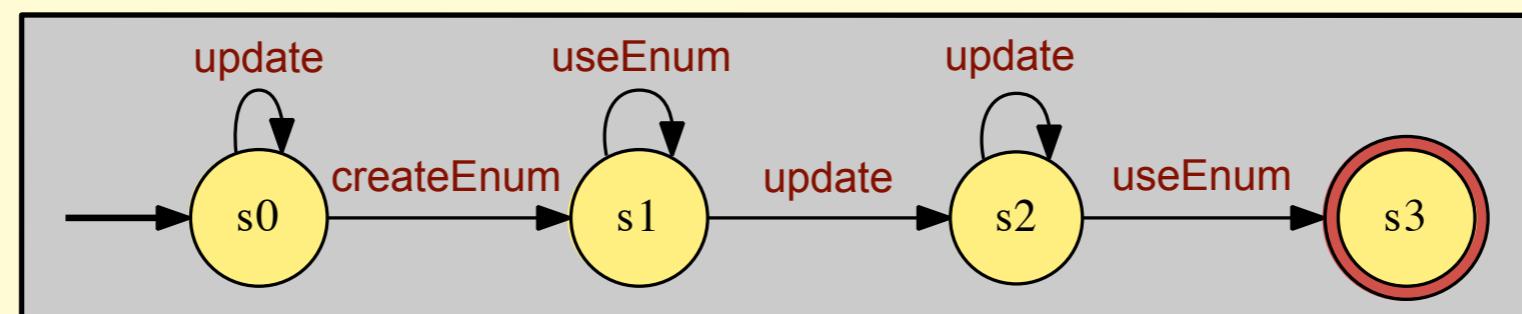
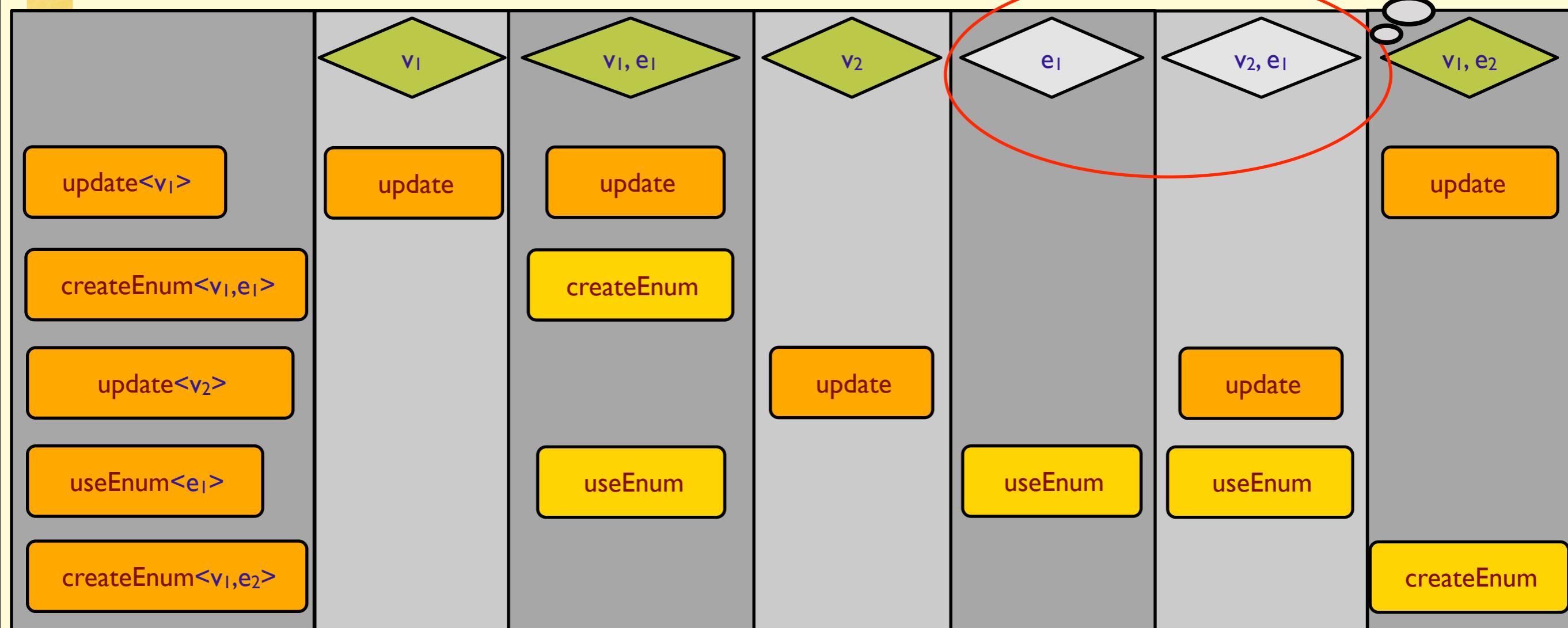


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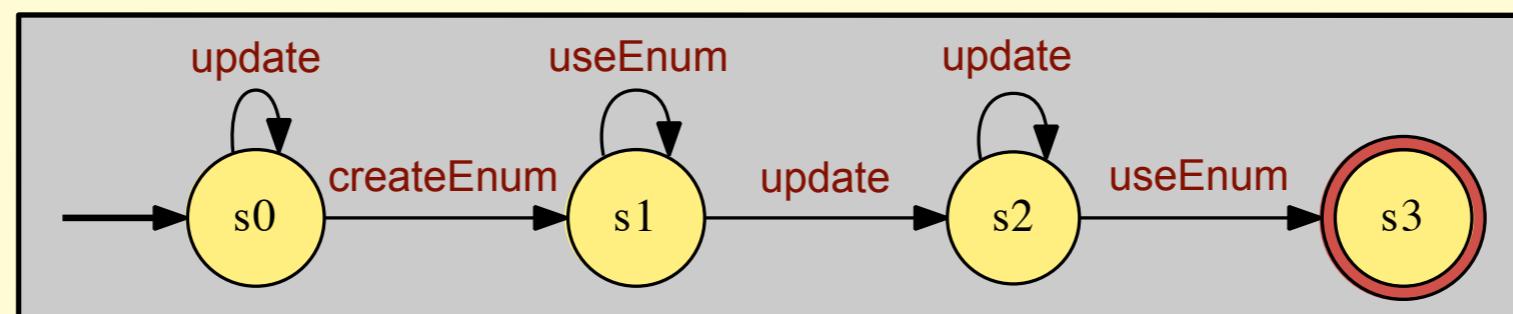
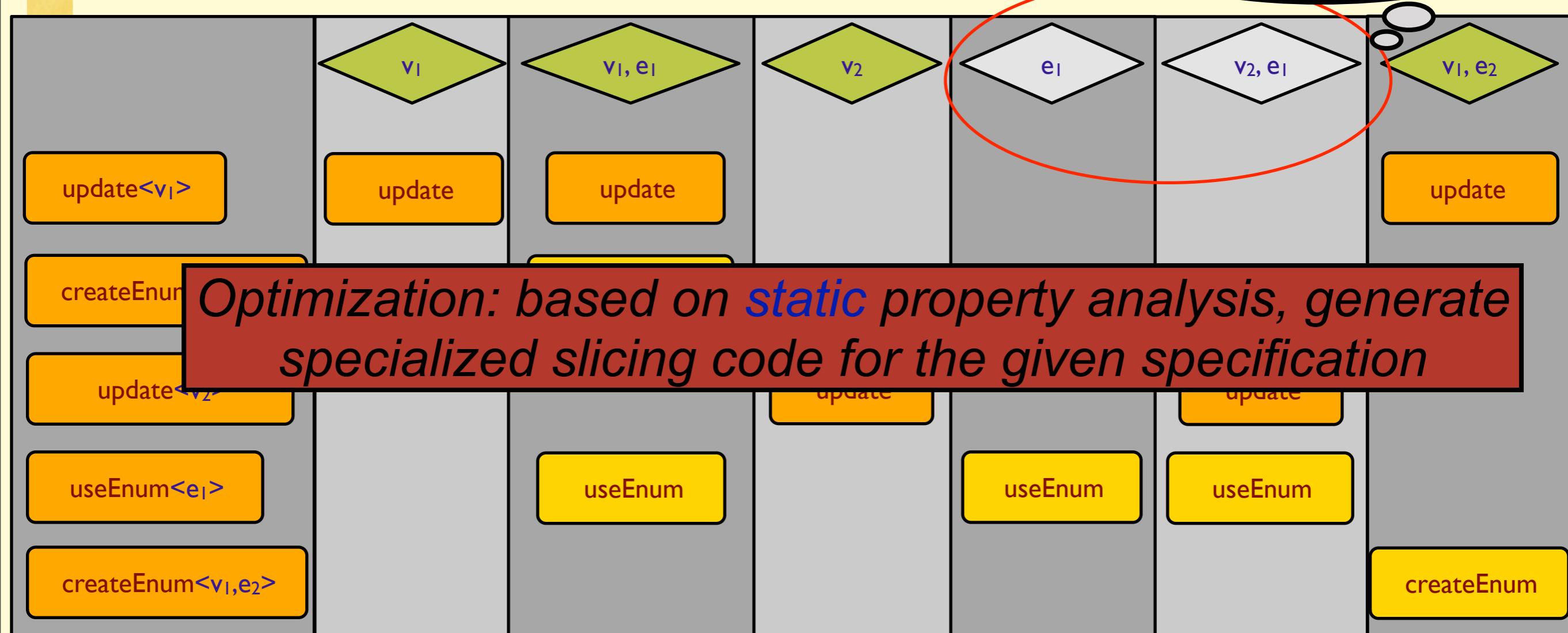
Example Run

Real programs generate a LOT of worthless monitors!



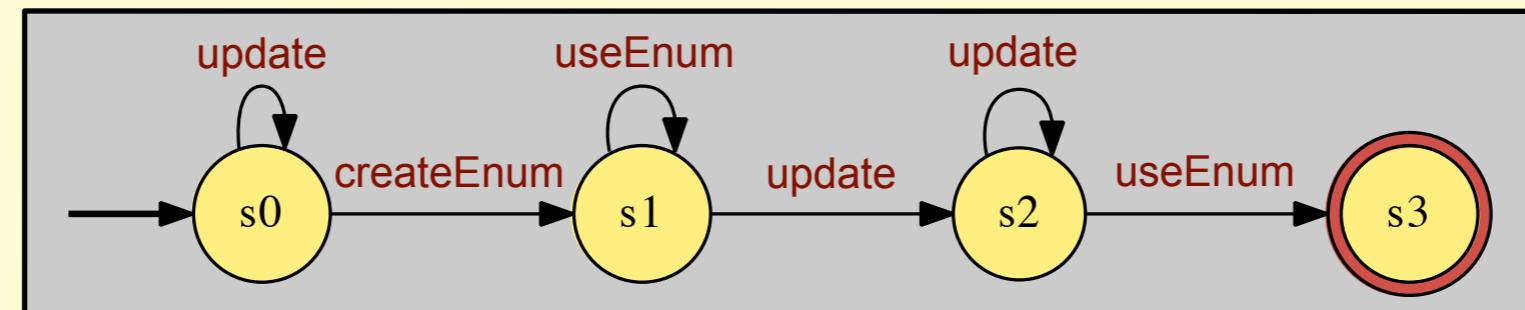
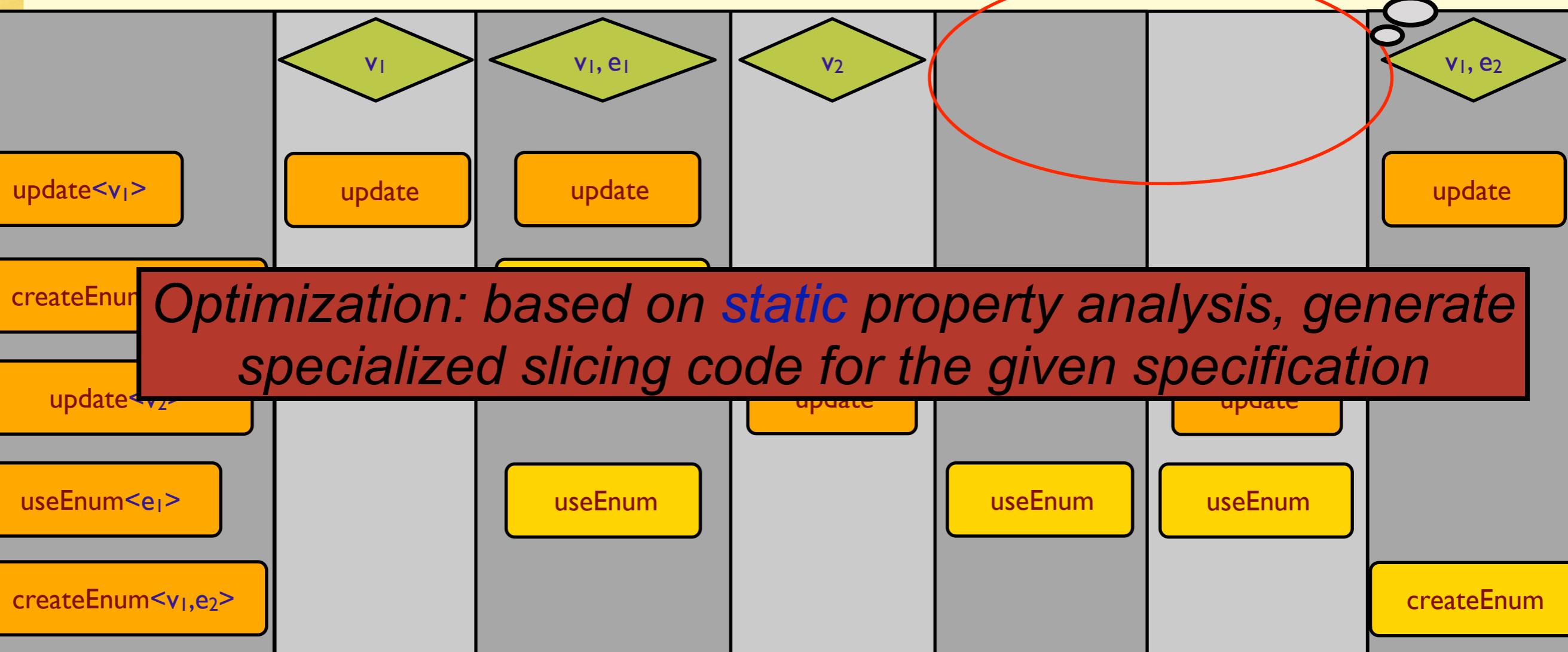
Example Run

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Example Run

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Enable Set Optimization

- We need to remove worthless monitors!
- Two directions for optimization:
 - Program analysis [AOSD '09]
 - Property analysis [Here]
- Enable Sets are a way to abstract knowledge of the property
 - Solves the problem of useless monitor generation
 - Improves performance radically in some cases without slowdowns in any cases

Enable Sets

- For each event which parameters must be instantiated before it may occur
- Computed by the monitor generation code for each formalism
 - Monitor generation code does not know about parameters
 - Compute in terms of events, MOP infers parameters

For trace t , $\text{set}(t) = \{e \mid e \in t\}$, $\text{enable}(e, L) = \{\text{set}(t_1) \mid t_1 et_2 \in L\}$

Enable Sets

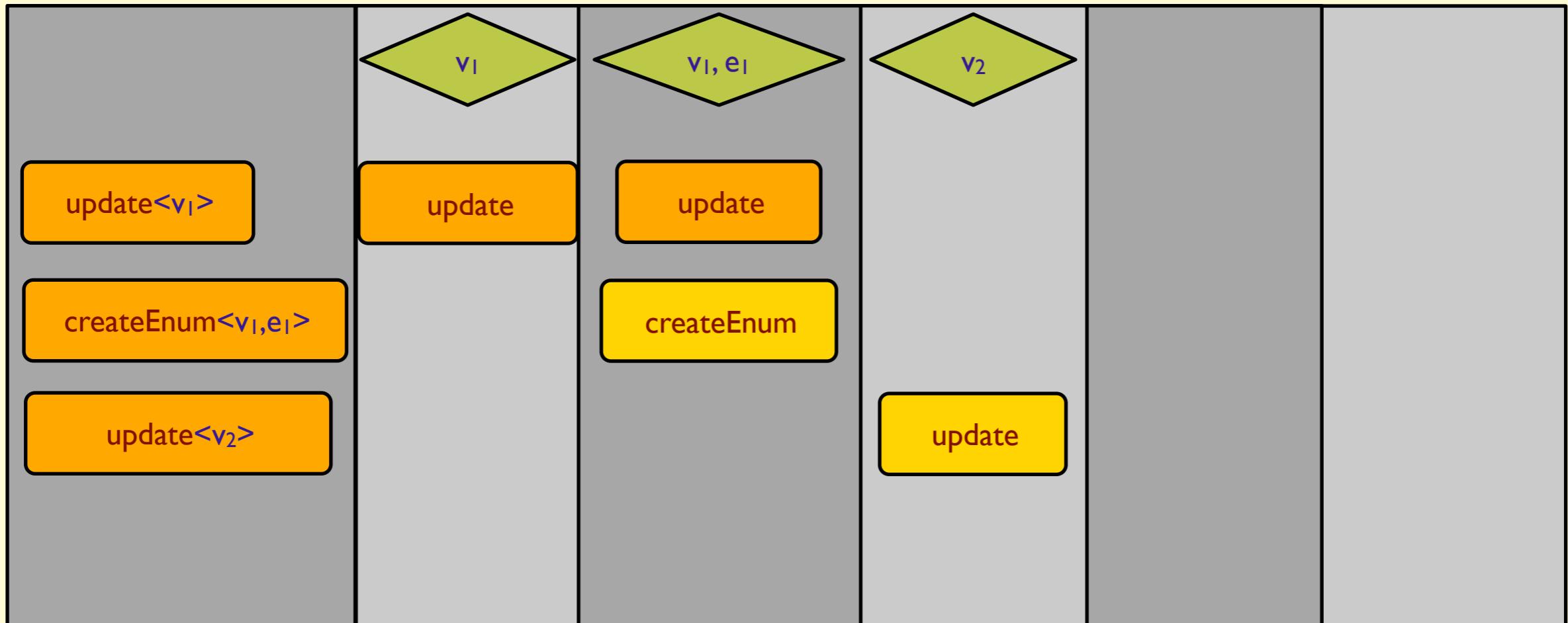
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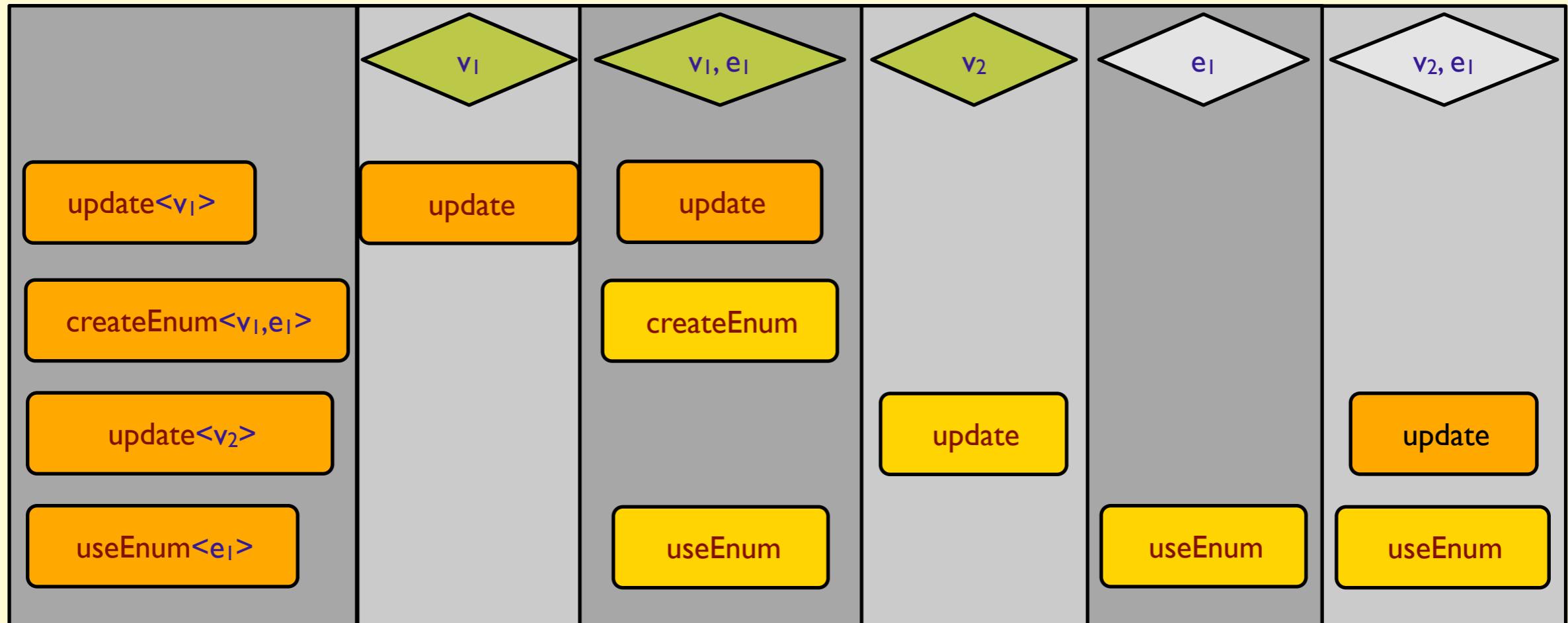
Possible traces - **createEnum useEnum update useEnum**
createEnum update useEnum

useEnum - **{ {createEnum}, {update, createEnum}, {update, createEnum, useEnum} }**

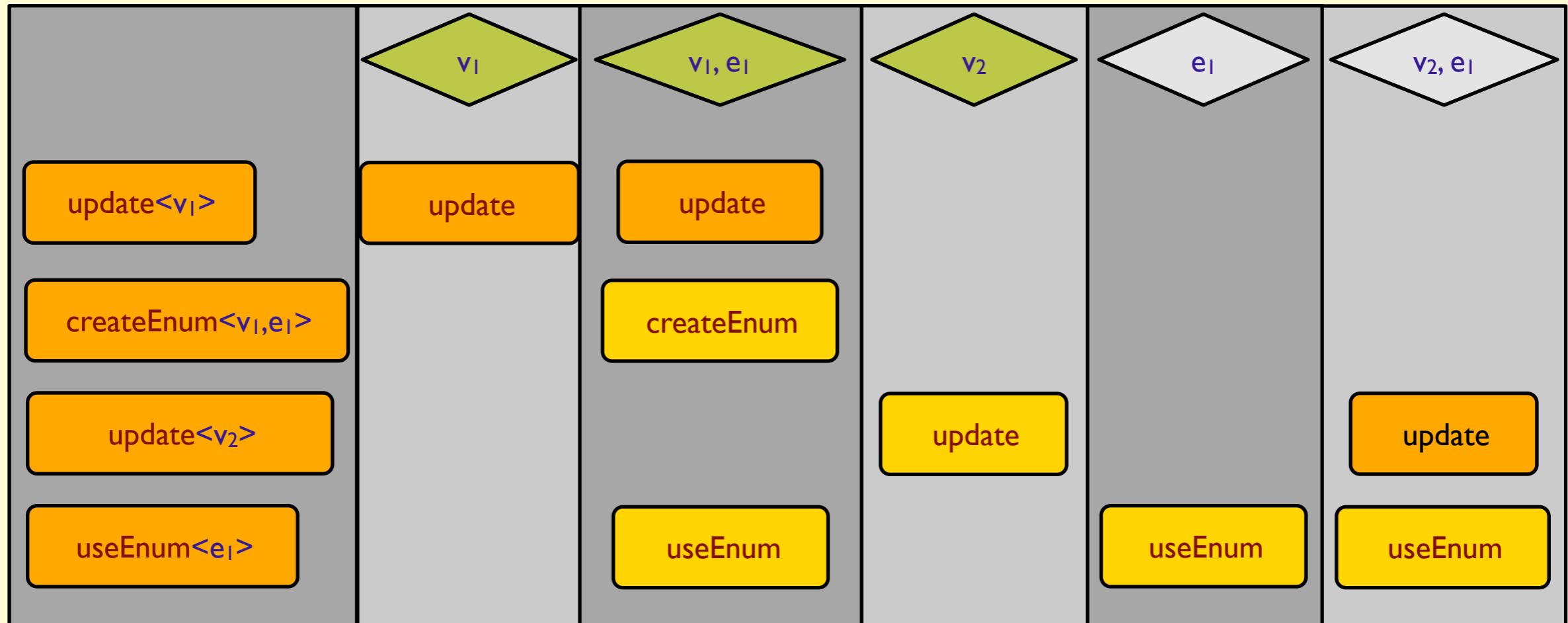
Optimizing with Enable Sets



Optimizing with Enable Sets

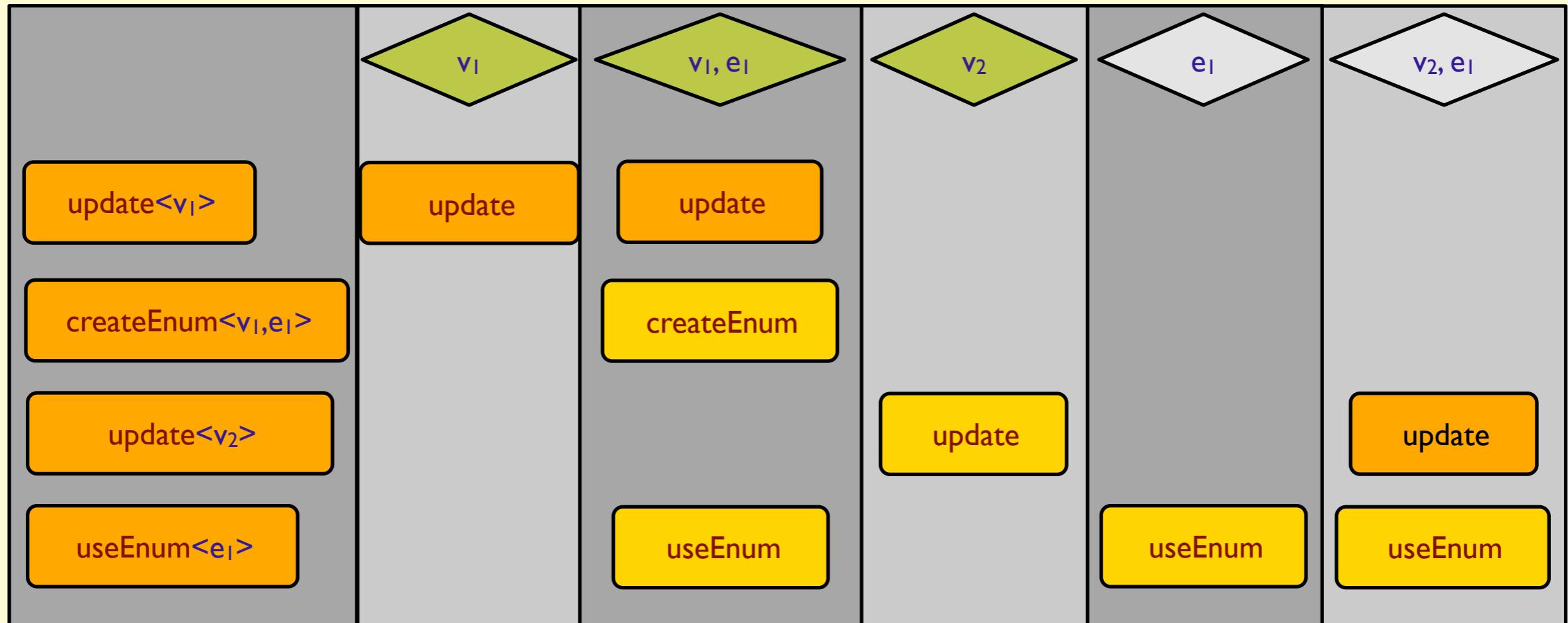


Optimizing with Enable Sets



`useEnum` - {{`createEnum<v,e>`},{`update<v>`, `createEnum<v,e>`},
{`update<v>`, `createEnum<v,e>`, `useEnum<e>`}}

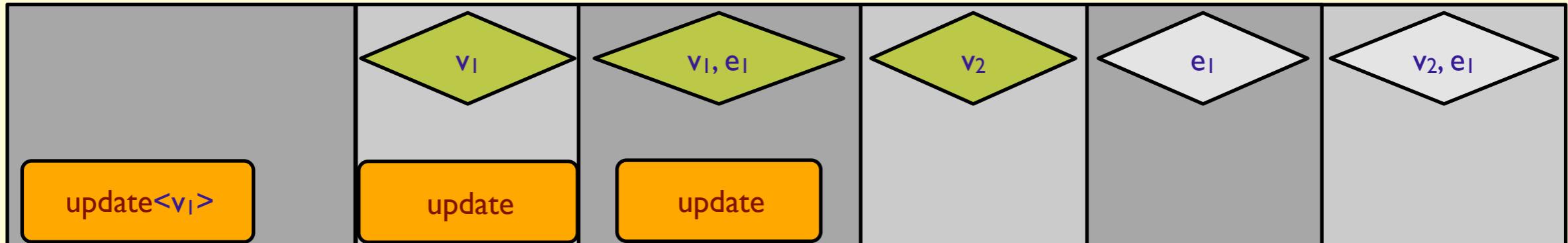
Optimizing with Enable Sets



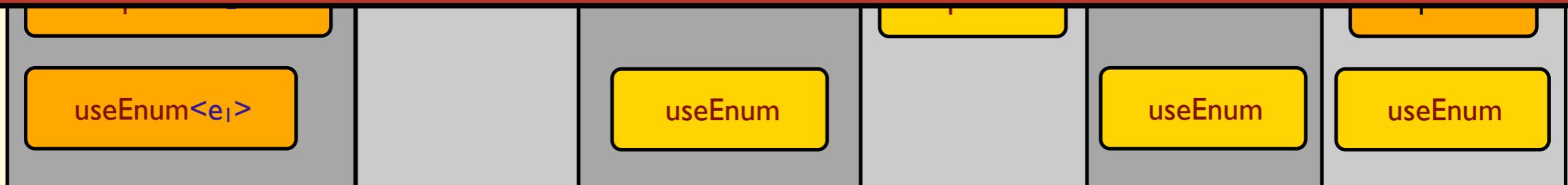
`useEnum` - {{`createEnum<v,e>`},{`update<v>`, `createEnum<v,e>`},
{`update<v>`, `createEnum<v,e>`, `useEnum<e>`}}

`useEnum` - {{ v, e }}

Optimizing with Enable Sets



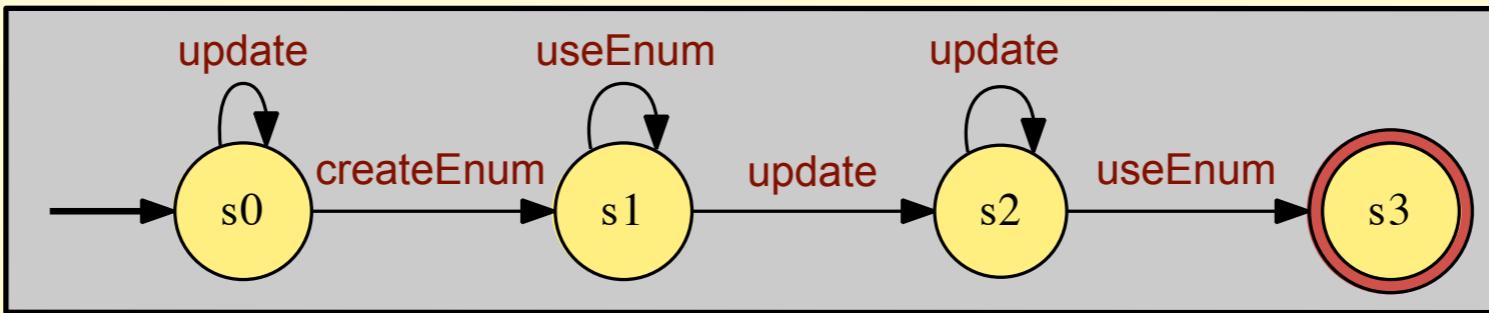
Because **v** and **e** are not both instantiated before **useEnum**, **<e₁>** and **<v₂,e₁>** will not be created!



useEnum - {{**createEnum<v,e>**}, {**update<v>**, **createEnum<v,e>**},
{**update<v>**, **createEnum<v,e>**, **useEnum<e>**}}

useEnum - {{**v,e**}}

Computing Match Enable Sets for Finite State Machines



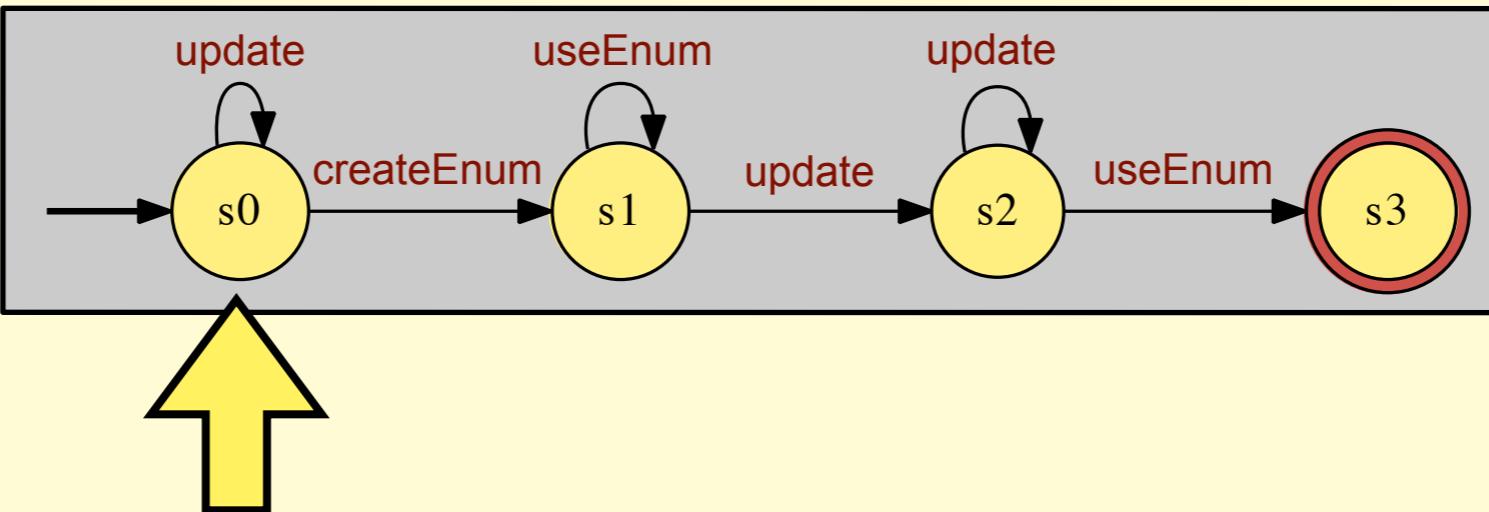
current set - {}

update - {}

createEnum - {}

useEnum - {}

Computing Match Enable Sets for Finite State Machines



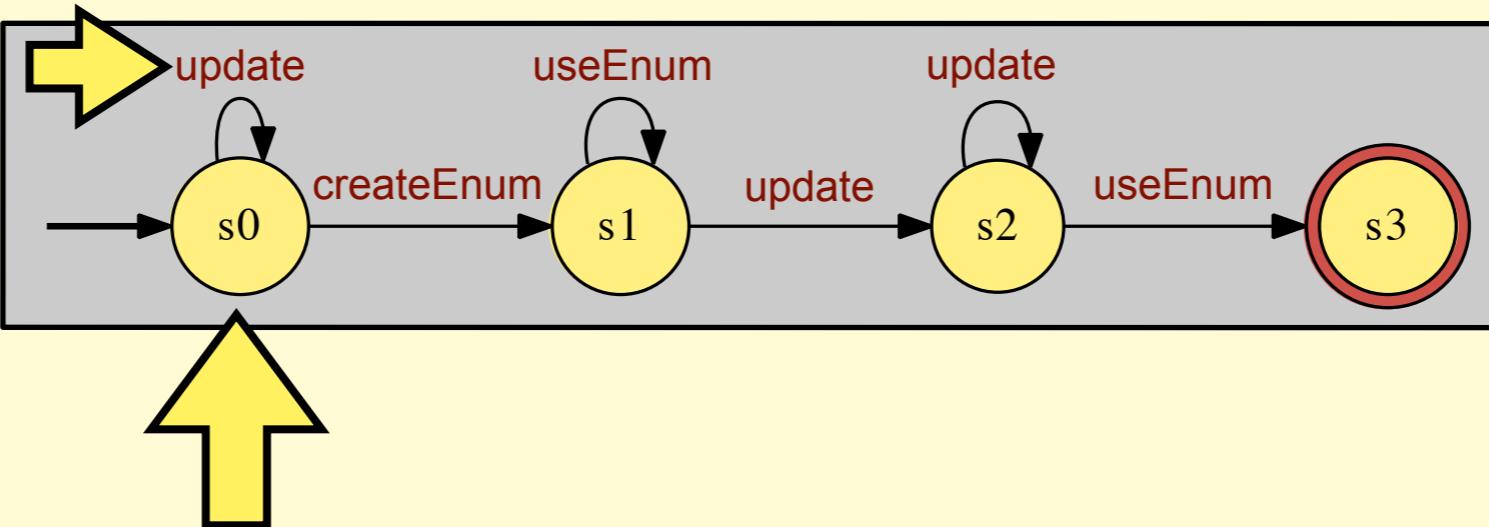
current set - {}

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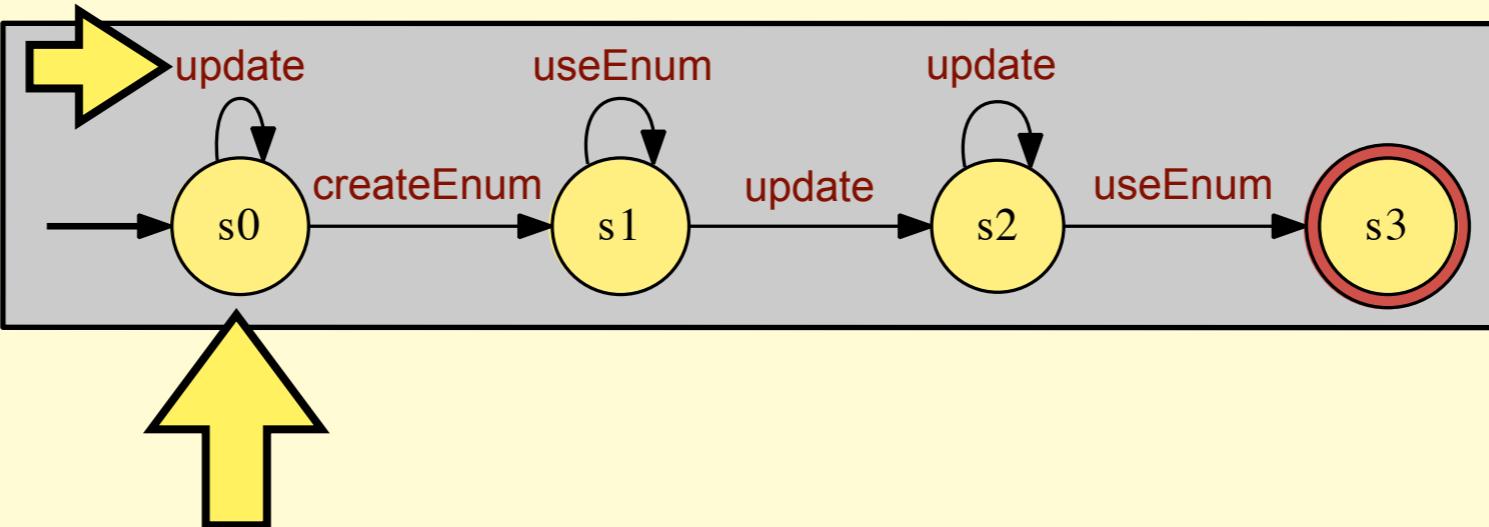
current set - {}

update - {}

createEnum - {}

useEnum - {}

Computing Match Enable Sets for Finite State Machines



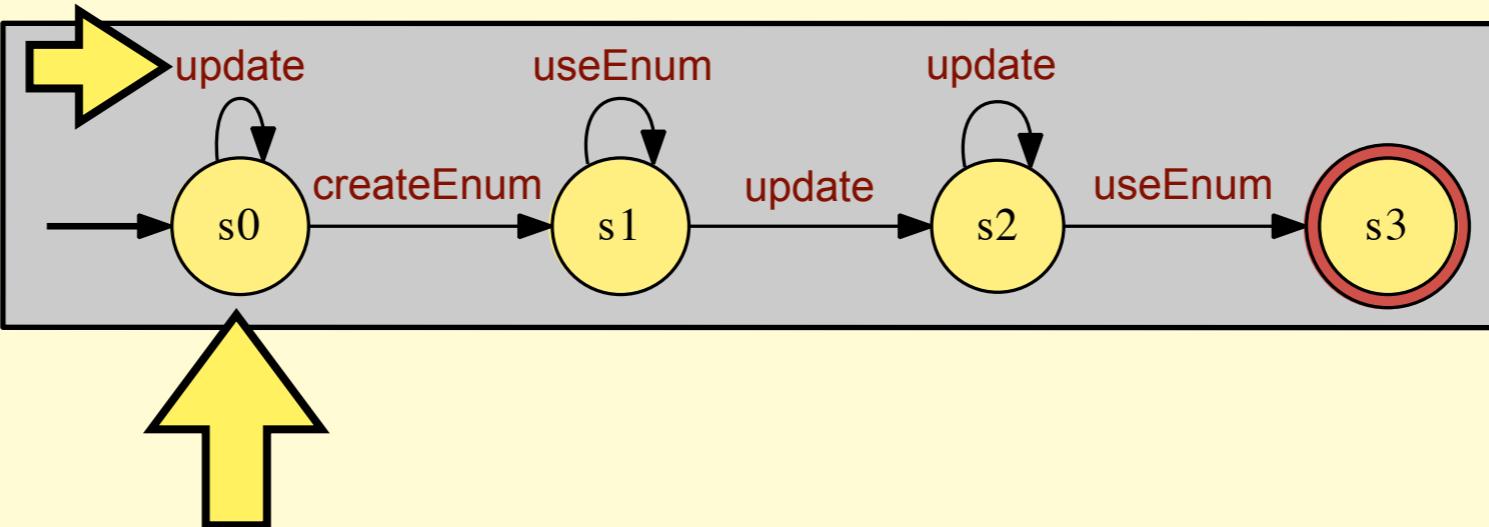
current set - {}

update - {{}}

createEnum - {}

useEnum - {}

Computing Match Enable Sets for Finite State Machines



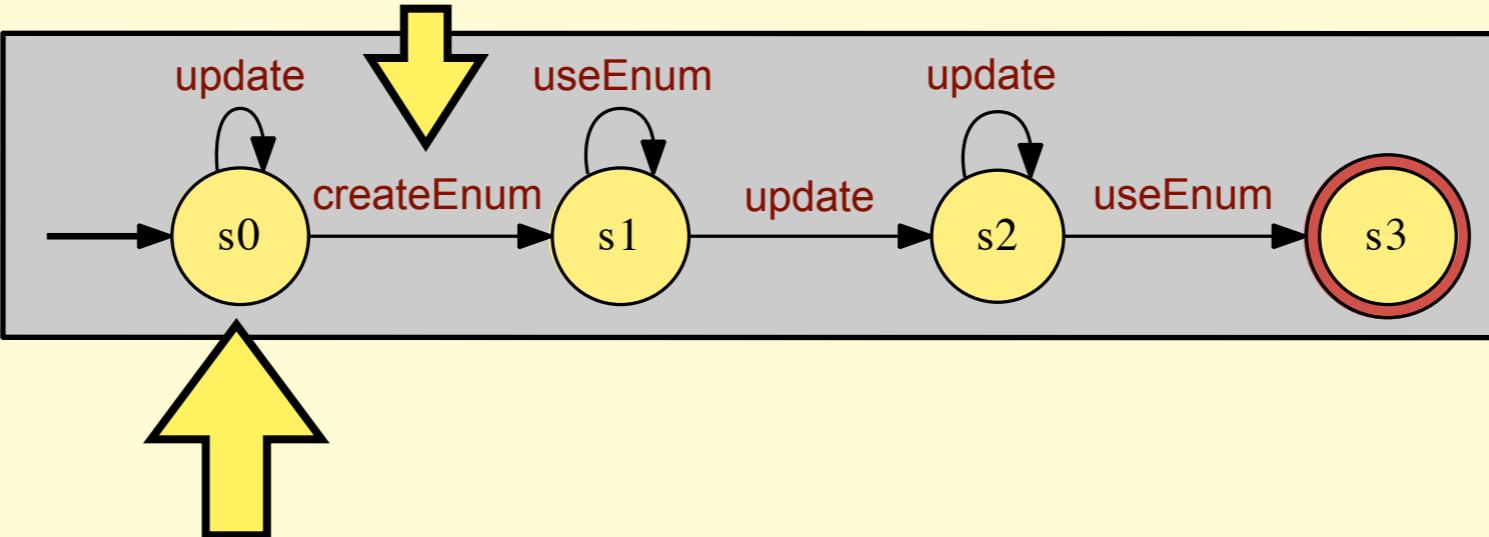
current set - {**update**}

update - {}

createEnum - {}

useEnum - {}

Computing Match Enable Sets for Finite State Machines



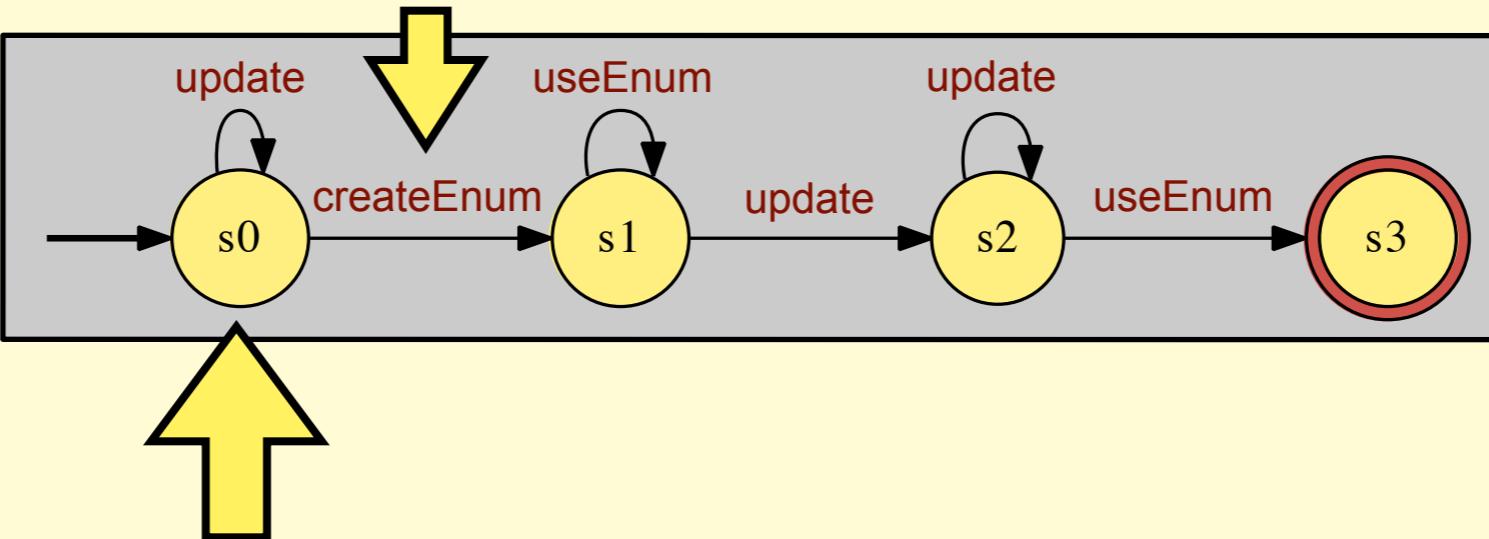
current set - {**update**}

update - {}

createEnum - {}

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Computing Match Enable Sets for Finite State Machines



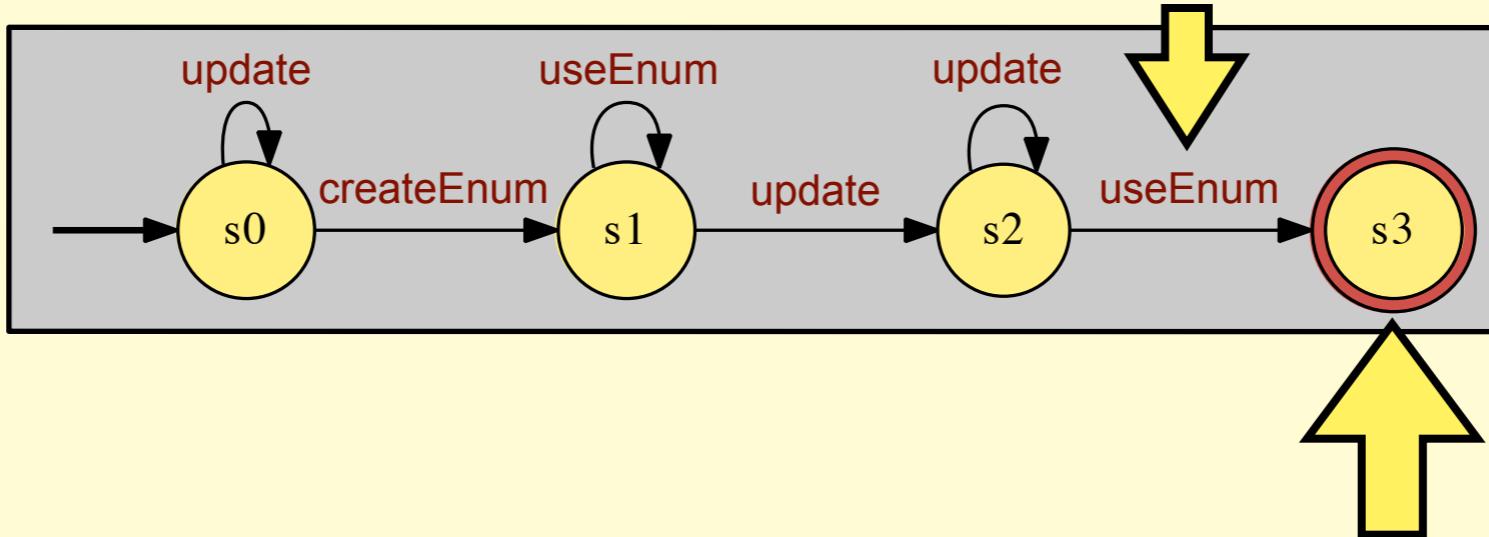
current set - {**update**}

update - {}

createEnum - {{**update**}}

useEnum - {}

Computing Match Enable Sets for Finite State Machines



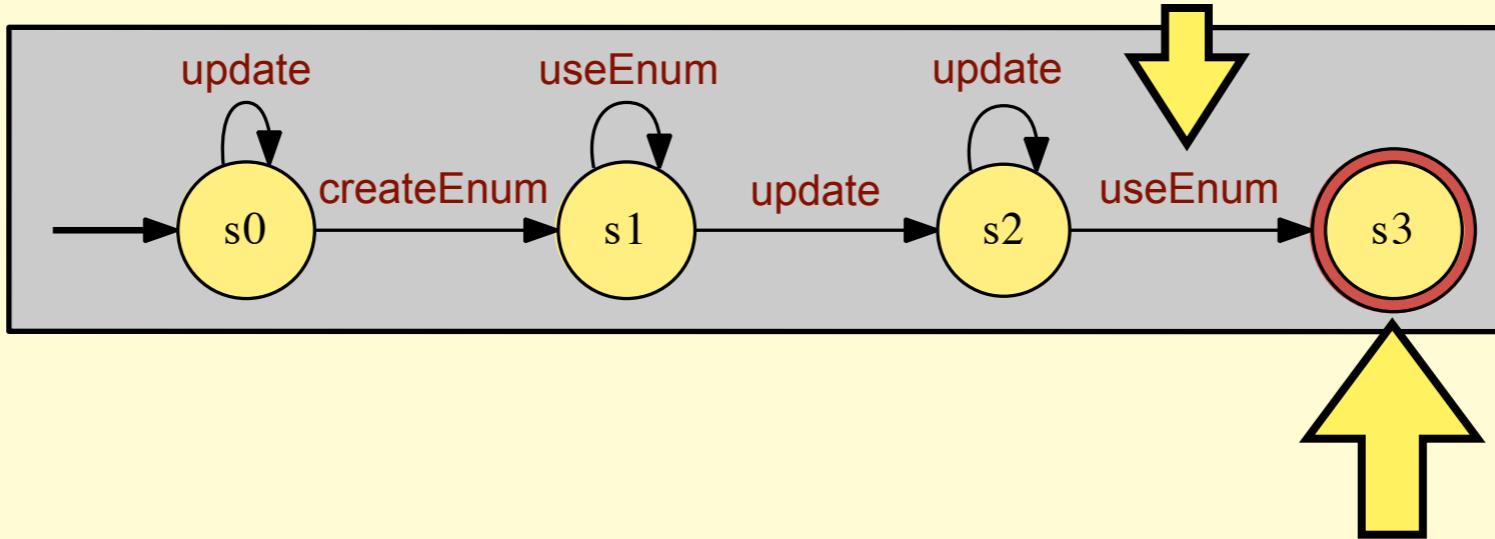
current set - {**update, createEnum, useEnum**}

update - {{}, {**update, createEnum, useEnum**}}

createEnum - {{**update**}}

useEnum - {{**update, createEnum**},
{{**update, createEnum, useEnum**}}}

Computing Match Enable Sets for Finite State Machines



current set - {**update, createEnum, useEnum**}

update - {{}, {**update, createEnum, useEnum**}}

createEnum - {{**update**}}

useEnum - {{**update, createEnum**},
{**update, createEnum, useEnum**}}

Repeated for each path that results in a different set

Computing Enable Sets for Context Free Grammars

- We also can compute enable sets for the matching context free grammars
- They can be computed as the least fixed point of these equations:

$$\text{GenSets}(\epsilon) = \{\emptyset\}$$

$$\text{GenSets}(e) = \{\{e\}\}$$

$$\text{GenSets}(A) = \bigcup_{A \rightarrow \gamma} \text{GenSets}(\gamma)$$

$$\text{GenSets}(\gamma_1 \gamma_2) = \{T \cup U \mid T \in \text{GenSets}(\gamma_1), U \in \text{GenSets}(\gamma_2)\}$$

$$\text{PreSymSets}(S) = \{T \cup U \mid A \rightarrow \gamma_1 S \gamma_2, T \in \text{PreSymSets}(A), U \in \text{GenSets}(\gamma_1)\}$$

$$\text{enable}(e, \text{match}) = \text{PreSymSets}(e)$$

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DaCapo - Benchmark Suite Used for Evaluation

- A popular benchmark suite for Java
 - www.dacapobench.org
- Benchmark forms of non-trivial programs
 - Bloat: A bytecode optimizer
 - Antlr: A parser generator
 - Eclipse: A popular Java IDE
 - Jython: Python for the Java Virtual Machine
 - ...

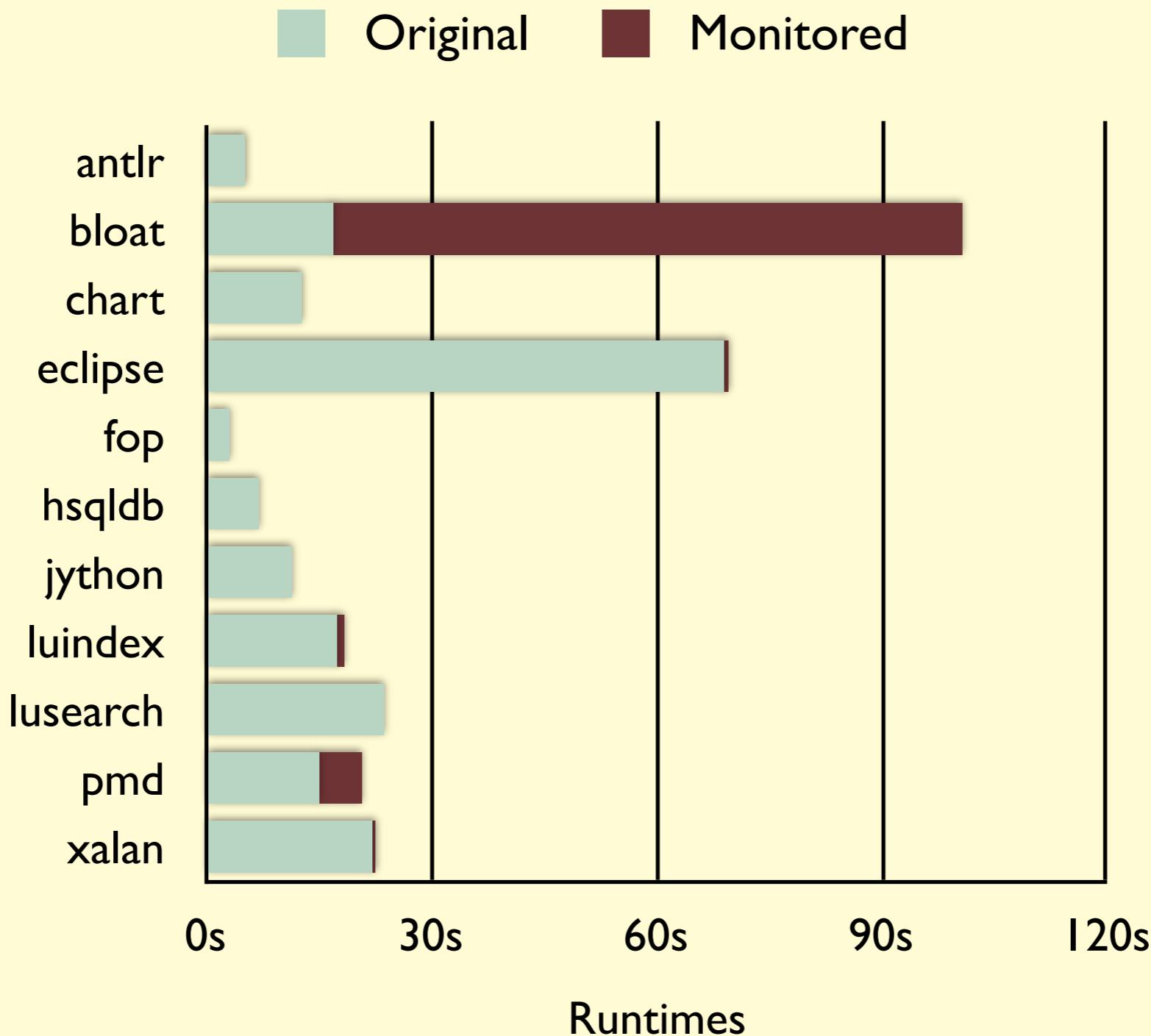
Properties Used for Evaluation

- **UnsafeMapIterator**
 - A Map m must not be updated while iterating over Collection c backed by m
 - $\text{createColl} < m, c >$ $\text{updateMap} < m >^*$ $\text{createlter} < c, i >$
 $\text{uselter} < j >^*$ $\text{updateMap} < m >^+$ $\text{uselter} < j >$
- **UnsafeSyncCollection**
 - A synchronized Collection c must not be accessed asynchronously
 - $(\text{sync} < c > \text{asyncCreatelter} < c, i >)$
| $(\text{sync} < c > \text{syncCreatelter} < c, i > \text{accesslter} < i >)$

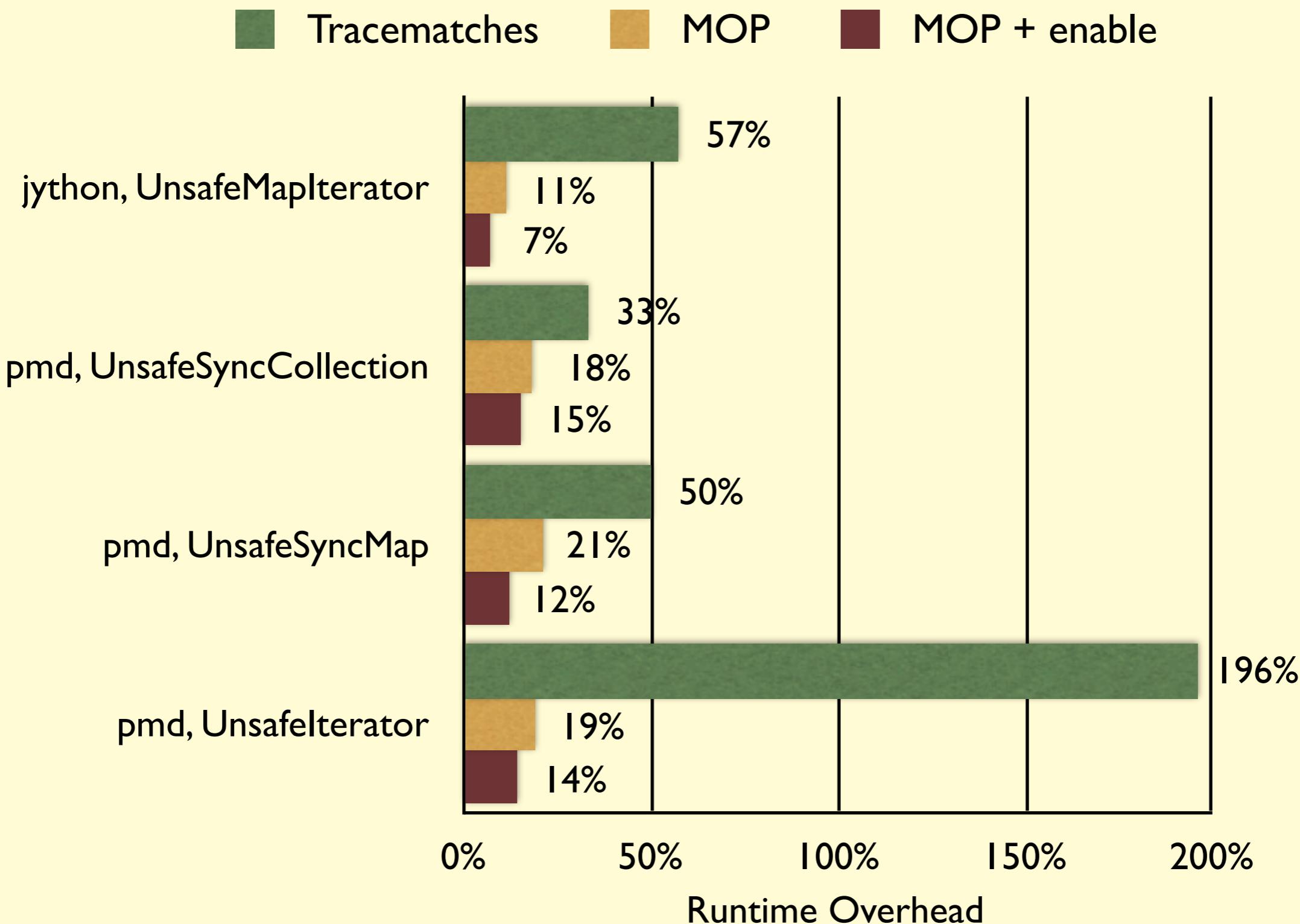
More Properties...

- **UnsafeSyncMap**
 - A Collection *c* created from a synchronized Map *m* must not be accessed in asynchronously
 - *sync<m> createSet<m,c>(asyncCreatelter<c,i>|(syncCreatelter<c,i> accesslter<i>))*
- **SafeFile**
 - A file opened in a given method must be closed in the same method
 - *S -> S begin<t> S end<t> | S open<f> A close<f> | ε*
A -> A begin<t> A end<t> | ε

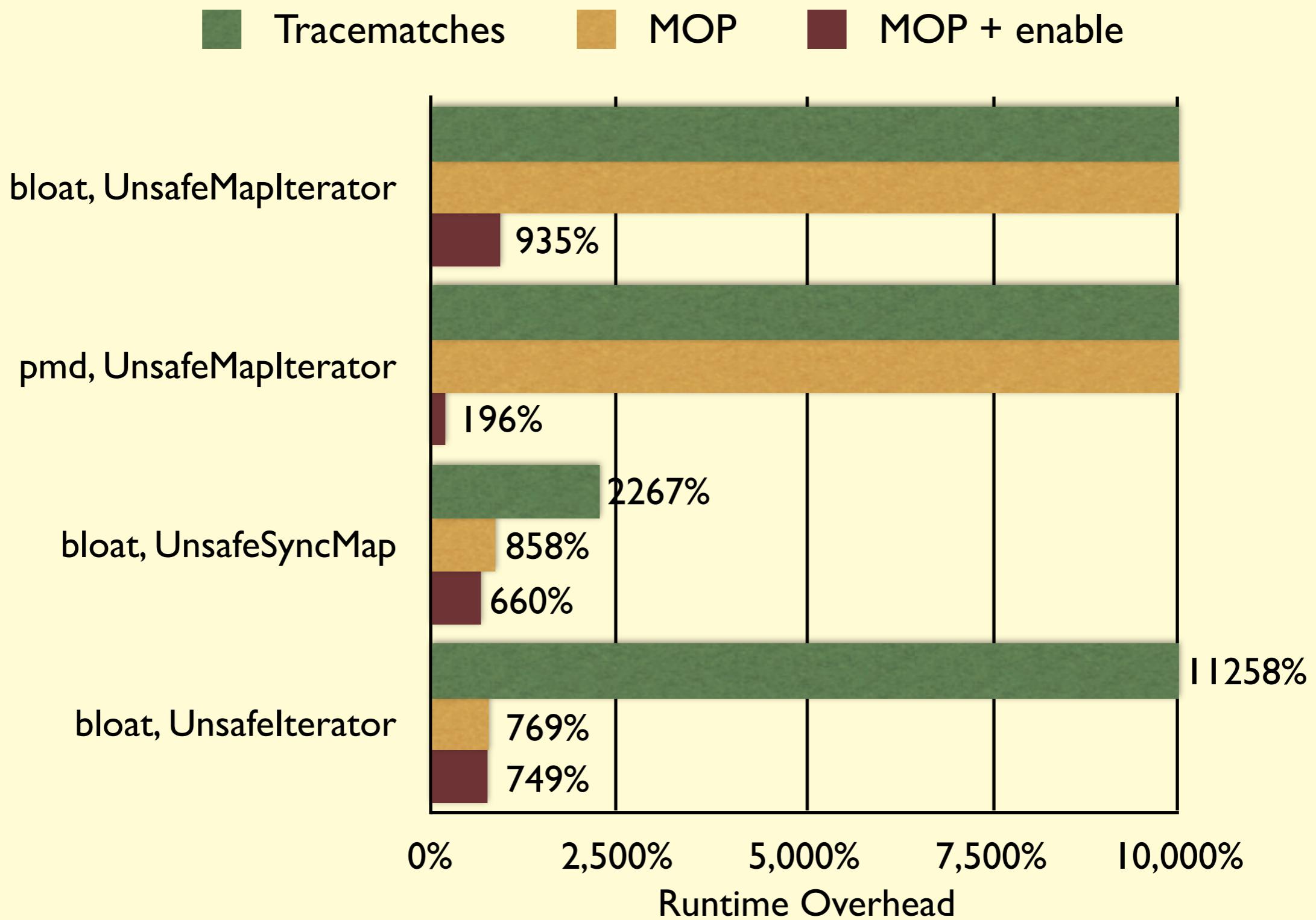
Original and Monitored Program Runtimes



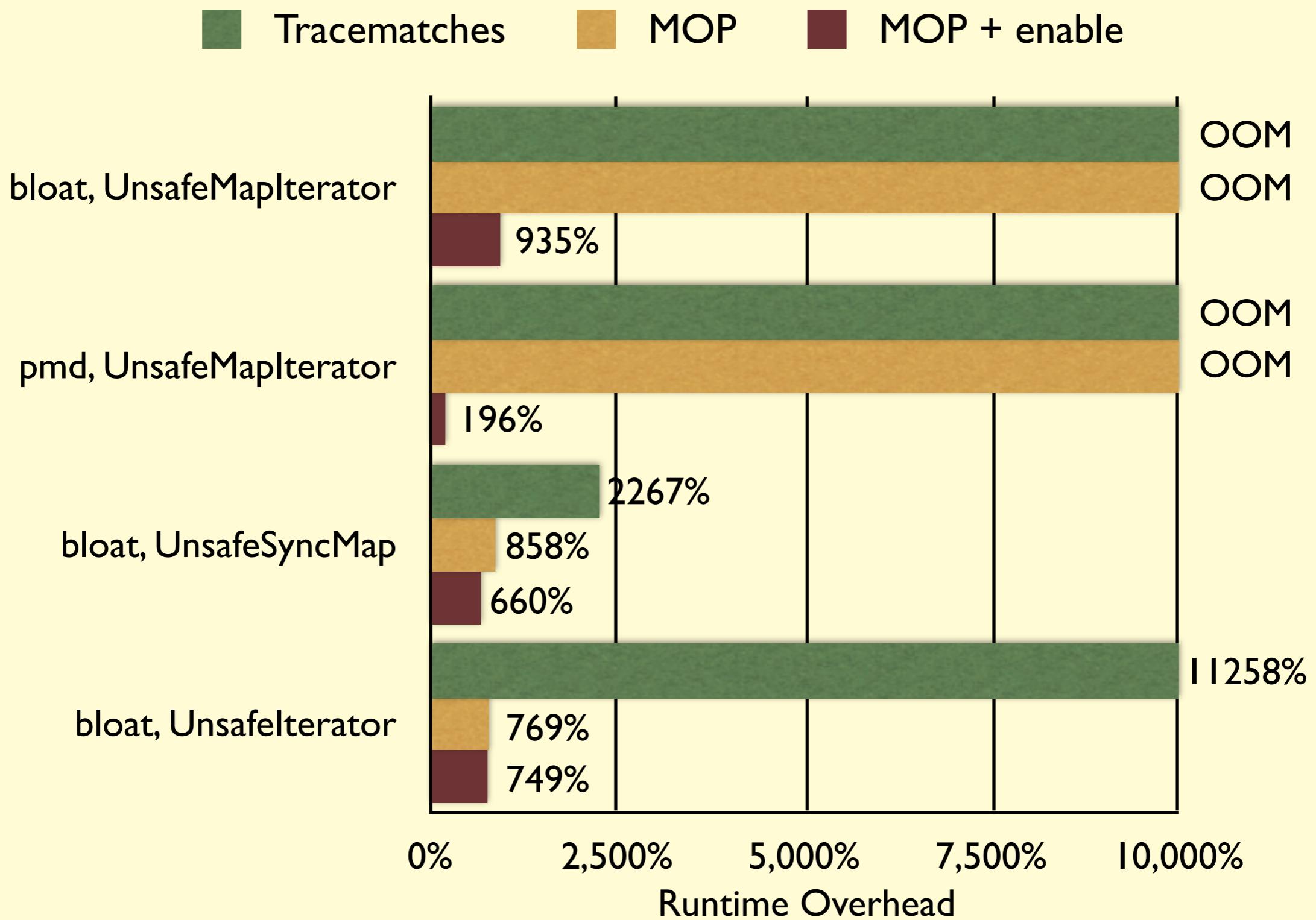
Non-Trivial Overheads



The Worst Overheads



The Worst Overheads



Conclusions and Future Work

- Parametric properties are useful and can be feasibly monitored
- Parametric trace slicing
 - Generic with respect to property formalism
- Enable set optimization makes trace slicing efficient
- Evaluation
 - Low runtime overhead
- Extend idea of Enable Sets to monitor termination